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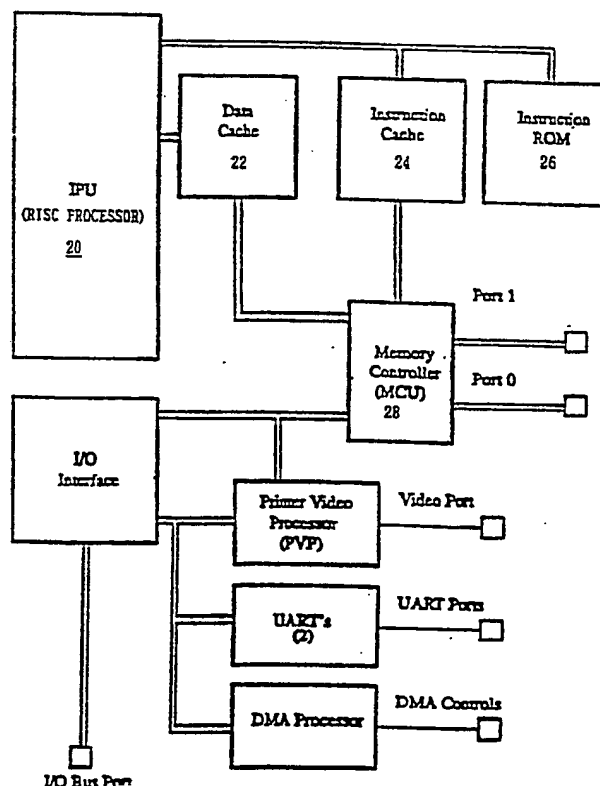
(51) International Patent Classification ⁵ : G06K 15/00, 15/02	A1	(11) International Publication Number: WO 93/01565 (43) International Publication Date: 21 January 1993 (21.01.93)
(21) International Application Number: PCT/JP92/00867 (22) International Filing Date: 7 July 1992 (07.07.92) (30) Priority data: 726,929 8 July 1991 (08.07.91) US (71) Applicant: SEIKO EPSON CORPORATION [JP/JP]; 4-1, Nishi-Shinjuku 2-chome, Shinjuku-ku, Tokyo 163 (JP). (72) Inventors: LENTZ, Derek, J. ; 17400 Phillips Avenue, Los Gatos, CA 95032 (US). WANG, Johannes ; 25 King Street, Redwood City, CA 94062 (US). DEOSARAN, Trevor ; 555 East Washington Avenue, Sunnyvale, CA 94086 (US). YOUNG, Linley, M. ; 2209 Calle Tiara, La Jolla, CA 92037 (US). YAP, Kian-Chin ; 1233 Hazelett Court, San Jose, CA 95131 (US). NGUYEN, Le Trong ; 15096 Danielle Place, Monte Sereno, CA 95030 (US). MATSUBAYASHI, Makoto ; Seiko Epson Corporation, 3-5, Owa 3-chome, Suwa-shi, Nagano-ken 392 (JP). LAU, Te-Li ; 411 College Avenue, Apt. E, Palo Alto, CA 94306 (US).		(74) Agents: SUZUKI, Kisaburo et al.; Seiko Epson Corpora- tion, 4-1, Nishi-Shinjuku 2-chome, Shinjuku-ku, Tokyo 163 (JP). (81) Designated States: JP, KR, European patent (AT, BE, CH; DE, DK, ES, FR, GB, GR, IT, LU, MC, NL, SE). Published With international search report.

(54) Title: SINGLE CHIP PAGE PRINTER CONTROLLER

(57) Abstract

A page printer controller comprising an instruction processor, a printer video processor having a video port, a data and instruction cache memory, a memory controller for interconnecting the printer video processor with external memory, and with the data and instruction cache memory, and an I/O controller for interconnecting the page printer controller to an I/O bus. The printer video processor is integrated with the I/O controller and the memory interface for optimizing memory bandwidth and allowing priority memory access. The instruction processor interfaces with a host computer through the I/O bus. The instruction processor includes pixel modification circuitry, orthogonal rotator circuitry, and bit/byte mirror circuitry. The printer video processor interfaces with the instruction processor through the I/O bus, and interfaces with memory through the memory interface. The printer video processor also controls a printer through the video port.

PPC Block Diagram



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DESCRIPTION

SINGLE CHIP PAGE PRINTER CONTROLLER

Cross-Reference to Related Application

This application is related to the following co-pending patent applications:

5 "High Performance RISC Microprocessor Architecture", Serial No. 07/727.006, filed July 8, 1991 (SMOS-7984MCF/GBR).

"Microprocessor Architecture Capable of Supporting Multiple Heterogeneous Processor", Serial No. 07/726.893, filed July 8, 1991 (SMOS-7992MCF/GBR/WMB).

Background of the Invention

10 This invention relates generally to controllers for printers such as laser printers, and more particularly the invention relates to a processor controller for a printer which can be implemented in a very large scale integrated circuit (VLSIC) semiconductor chip.

15 A page printer controller operates under control of a host computer in providing the video signals for a printer such as a laser printer. The functions of transforming data to a display coordinate system, clipping objects to conform to an image window, and scaling the window to the viewpoint of the
20 destination display device are discussed in U.S. Patent No. 4,449,201 for "Geometric Processing System Utilizing Multiple Identical Processors" and in U.S. Patent No. 4,648,045 for "High Speed Memory and Processor System for Raster Display," as applied to graphics display.

Disclosed in the above-referenced co-pending patent applications are a unique single chip RISC processor and a memory control unit operable therewith to permit multiprocessing capability. The present invention is directed to a printer controller architecture which can utilize the RISC processor architecture and interface with the memory control unit. The controller includes peripheral components such as a printer video processor (PVP), UART ports, an I/O bus interface, a DMA controller, a firmware ROM, and memory interface. The major functions of the printer controller are to communicate with a host computer, transform data into image data (coordinate transformation, clipping, scaling, and rasterization) interface with other components of the print engine (e.g. laser printer), and run diagnostics. Known in the prior art is a single chip unit (Intel i860) including a RISC processor and special graphics hardware. However, the graphics hardware is outside of the RISC processor and not physically and functionally integrated therein.

Summary of the Invention-

Accordingly, an object of the invention is a printer controller which can be implemented in a single VLSIC semiconductor chip.

Another object of the invention is a printer controller which can utilize a RISC microprocessor architecture as disclosed in the above co-pending application.

A feature of the invention is the provision of special function circuits within a RISC microprocessor which are used by special graphics functions and integrated into the processor in the same way as other function units.

Another feature of the invention is a printer video processor having efficient and integrated direct memory access.

The invention and objects and features thereof will be more readily apparent from the following detailed description and appended claims when taken with the drawings.

Brief Description of the Drawing-

FIGURE 1 illustrates the PPC system block diagram.

FIGURE 2 illustrates the PPC block diagram.

FIGURE 3 illustrates pixel modification datapath.

5 FIGURE 4 illustrates the orthogonal rotator.

FIGURE 5 illustrates bit/byte mirror.

FIGURES 6A and 6B illustrate the integration of the special function circuitry of FIGURES 3-5 into the RISC processor datapath.

10 FIGURE 7 illustrates video processor block diagram.

FIGURE 8 illustrates memory system structure.

FIGURE 9 illustrates data chip, two-way interleaved.

FIGURE 10 illustrates address chip, two-way interleaved.

FIGURE 11 illustrates horizontal timing.

15 FIGURE 12 illustrates vertical timing.

FIGURE 13 illustrates UART O output timing.

FIGURE 14 illustrates UART O input timing.

FIGURE 15 is a functional diagram of the I/O controller datapath.

20 FIGURES 16-21 are timing diagrams illustrating I/O bus transactions.

FIGURE 22 illustrates DMA read cycles.

FIGURE 23 illustrates DMA write cycles.

Detailed Description of Illustrative Embodiments-

25 The page printer controller (PPC) in accordance with the invention can be implemented using the RISC processor and memory controller described in the copending applications, and the printer controller described herein for controlling a page printer engine is based thereon. Reference to "laser" printers
30 may be generalized to refer to any page type printer with a similar interface.

Referring now to the drawings, FIG. 1 illustrates the general system block diagram of the PPC 10 as it interfaces with a two-way interleaved memory array 12 including memory

bank 0 and optional memory bank 1, an I/O bus 14 for communicating with a host computer and various peripherals; a special purpose (printer control) UART, a general purpose UART, and providing the video signals to the raster scan laser of the printer. The PPC unit 10 is further illustrated in the block diagram of FIG. 2 and includes a RISC CPU or instruction processing unit (IPU) 20 which communicates directly with an optional data cache 22, an optional instruction cache 24, and an optional RISC instruction ROM 26. The microprocessor unit 20 further communicates with memory through the memory interface unit 28 having one or more memory ports.

As noted above, the microprocessor unit 20 in the preferred embodiment is based on the RISC chip disclosed in co-pending application Serial No. 07/727.006 SMOS-7984MCF/GBR, and the memory interface disclosed in co-pending application Serial No. 07/726.893, SMOS-7992MCF/GBR/WMB, both of which are incorporated by reference. In addition the instruction processing unit has several special functional units in support of the printer video processor 30 including a pixel modification datapath, an orthogonal rotator, and a bit/byte mirror.

FIG. 3 is a functional block diagram of the pixel modification datapath in which pixels are modified by graphics instructions. The pixels are modified by graphics instructions as shown in the datapath. The destination and source are bit-wise combined by the LFU. The combined results are passed into the output MUX and either the original destination or the modified destination is passed to the output. The computations are bit-wise, so that the selection of the modified or unmodified destination is made by the LEM, REM and Mask Data of the individual datapath bit. The LEM and REM are 5 bit numbers which define the left and right edge of the bit range to be modified, and are processed by edge generator logic to expand into bit fields (the left mask and right mask) which are used to control the multiplexer. The left mask, the right

mask and the Mask Data are bit-wise logically ORed, such that all bits must be set to 0 for the destination bit to be changed. The circuit of FIG. 3 is a "bit slice" representation of the graphics pixel modification datapath. The LFU function code is held in a processor state register.

The destination is always combined with a "source". The source may be predetermined data or copied from a memory bit map. The source and destination may be combined with any of the 16 possible combinations of 2 binary bits. When a mask is used, the mask restricts the pixels modified to those within the mask image which are set.

FIG. 4 is a functional block diagram of the orthogonal rotator which is used to assemble blocks of orthogonally (90 or 270 degree) rotated bitmap data. The low byte of the source data is written to a bit of the 8 registers held in the unit. Up to 32 bytes are written to fill up the internal words and the orthogonally rotated data is then written to memory or used by the processor. If the words are read out from low to high register address, the rotation is 90 degrees, while reading the registers out from high to low register address rotates the data by 270 degrees.

FIG. 4 shows the basic operation of the logic. The source is loaded with 8 bits. Bit 0 is passed to all 32 bits of register 0. Bit 'n' is passed to all 32 bits of register 'n'. The bit selected is used by the bit decode to select the bit within the 8 words to be modified. In this way, one bit of each register is always loaded by one bit from the source. The data in each word is available for reading at any time. In addition, all words may be pre-loaded directly through "move" instructions and all registers may be cleared as a block, simultaneously, for initialization purposes.

The orthogonal rotator instructions consist of: move to and from A or B registers, preset rotator and rotate byte orthogonally. Preset orthogonal rotator sets the values of all bits in the orthogonal registers with one instruction. The

instruction has one parameter "level". This 8 bit parameter loads 1 bit into all bits of each of the 8 orthogonal registers. Bit <0> is loaded into all bits of orthogonal register 0, bit <n> is loaded into all bits of register n, etc. Rotate byte orthogonally performs the orthogonal rotation, described above on the low order byte of data from a general register. Rotate byte orthogonally has 2 parameters: the register holding the byte of data and the register holding the bit selected. The data from the registers are truncated to 8 and 5 bits, respectively.

FIG. 5 illustrates the bit/byte mirror which consists of four multiplexers (MUX) each of which receives eight bit bytes and outputs the bits in either the same order or in reverse order (mirrored) in accordance with the following table:

	Input Bit:	Output Bit:	Input Bit:	Output Bit
15	0	7	8	15
	1	6	9	14
	2	5	10	13
	3	4	11	12
20	4	3	12	11
	5	2	13	10
	6	1	14	9
	7	0	15	8
25	16	23	24	31
	17	22	25	30
	18	21	26	29
	19	20	27	28
	20	19	28	27
	21	18	29	26
30	22	17	30	25
	23	16	31	24

After the bits have been mirrored within a word, the word can be written back to memory, one byte at a time, or the bytes can be reordered using a barrel shifter circuit. The bit mirror instruction reads the source word from a register and writes the mirrored contents back to the destination register.

5 The mirroring and orthogonal rotation of bitmaps may be performed with hardware support. Two hardware functional units are provided which allow firmware and software to perform time consuming tasks relatively quickly. The bit-mirror unit mirrors all of the bits within each byte of the word loaded into it. The bytes may then be rearranged, if desired, by the use of shift instructions such as extract byte and merge byte. 180 degree rotations may be performed by this unit in conjunction with the standard shifter functions.

10 The orthogonal rotator unit places the low order byte of a word in a selectable bit of registers. One bit is placed in each register. The effect is that the contents of the 8 registers ends up being a 32 bit x 8 word block of bits which have been orthogonally rotated. The bit-mirror unit is provided to support the mirroring of bits within memory, for 15 orthogonal rotations of 180°. The unit swaps the bits within all bytes of the word passed through it. For convenience the bit-mirror unit is integrated into the pixel modification datapath.

20 FIGS. 6A and 6B illustrate the integration of the special function circuitry of FIGS. 3-5 into the IPU. The special function circuits may be integrated as function units or within function unit groups. In FIG. 6A, the circuitry is identified as a Function Unit (Group) and is interconnected with register files A and B through MUXes just as the other Function Units (e.g. ALU, Multiplier, FALU) are interconnected. FIG. 6B 25 further illustrates the interconnection of the source data MUXes and the control signal MUXes with the orthogonal Rotator, Pixel Modifier, and Mirror.

30 The graphics function units are completely integrated into the RISC IPU. The function units are used by special graphics instructions, which are executed in exactly the same way as other RISC instructions. The RISC core is capable of "super scaler" operation, where more than one instruction may be 35 executed and/or completed per cycle. The graphics function

units are scheduled by the schedule in the same way as the other, more standard, function units, such as the ALU's or multipliers. The orthogonal rotator and pixel modification units contain internal "state" information, which must be maintained during IPU context switches. These state registers are handled in the same way as other processor state information (e.g. divider quotient and remainder registers) through special register move instructions between the register files and the function units. The function units are available to the A and B IPU register files symmetrically through the same set of buses as the shifter function unit.

Printer Video Processor-

The printer video processor (PVP) 28 of FIG. 2 is further illustrated in the functional block diagram of FIG. 7. The PVP generates timing signals, performs scan data fetches and scan data serialization, and provides print engine synchronization as well as scan status and interrupt generation. Most of the circuitry is clocked by the system clock; however, the serializer and some control logic are clocked by another clock known as the print engine video clock (Wclk).

The video register file 30 contains registers used to generate timing and memory addresses. The registers may be read and written by the main system processor. The following registers reside in the video register file:

25	Register:	Purpose:
	Base	Frame buffer starting address (byte address)
	PBase	Previous base register (band buffer mode)
30	X_max	Width of frame buffer scan, in bytes
	Y_step	Memory address increment (see discussion), in bytes
	Y_max	Number of scans in frame
	T_margin	Number of scans at the top of the frame
	Bsize	Band buffer size in lines
35	V_int	Vertical interrupt scan number
	Base_current	Current base register (internal use)

	Current	Current scan memory address (internal use)
	X_current	Current X address within scan (internal use)
5	Y_current	Current scan number (internal use)
	Data_size	Memory access size (internal use)
	Bcnt	Current position in current band buffer
	Tmp_0	Temporary register (internal use)
	Tmp_1	Temporary register (internal use)
	Tmp_2	Temporary register (internal use)

10 The following registers are not in the register file, but are used for timing computations.

Register:	Purpose:
L_margin	Video clocks to the left of the first dot
X_max1	Program with (# of bytes in scan line -1)

15 The base memory address should be aligned to 16 byte boundaries. The scans may be odd numbers of bytes wide. X_max holds the number of bytes in a scan line rounded up to the nearest multiple of 8. X_max1 holds 1 less than the number of bytes per scan.

20 The constants ROM 32 are: 0, 1, 4, 8, 16.

The adder 34 performs arithmetic operations with the register file and constants. The results of computations, which may be written back into the register file, are used by the controller 35 to generate addresses, control and status.

25 The output may be passed to the memory address generator which then fetches the next word to be serialized. The state is set by the controller and the adder flags (zero and carry out).

The video data is loaded into a FIFO or Video Data Queue 36. The serializer is loaded from the queue when it is empty.

30 The queue has full and empty signals for the sequencer to thereby control the flow of data to prevent the loss of data or the emptying of the queue prior to the end of the scan. The depth of the queue is implementation dependent and is set by the expected worst case latency in the memory system and the

35 intended maximum video clock rate, but is a minimum of 8 words.

The serializer 38 is a buffered parallel-to-serial shift register, which generates data at 1 bit/dot from 8 bits of data. The shift register shifts the serial data out of the video interface serial data output pin (Wdata). The data shifts when the sequencer indicates data is required by the print engine, at the rate of 1 bit per video clock (DClk). DClk is equal to 1/8 or 1/16 of the Wclk in Canon mode. In front page printing, the lowest significant byte from the queue is loaded first into the shift register, whereas for the back page, the most significant byte is loaded first. In front page printing the data is shifted right, whereas it is shifted left for back page printing.

The use of the PPC in color printers requires the use of multiple passes to display multiple colors or the use of packed pixels in the frame buffer which are serialized out of the PPC at a multiple of the video clock required by the engine (e.g. 4x for 4 bits/pixel).

The mode, command and status register 40 may be read or written to by the main processor and provide software control of scan conversion, software monitoring of scan conversion, and software control of PVP interrupts.

The controller sequences the scan unit under the control of the command register (accessed by IPU 20). The IPU loads a scan command by setting the start frame bit of the command register, which tells the controller to start scanning. The controller initializes the scan and sequences through the page. The scan pending bit is reset when the Base register is free to be reloaded so that another scan may be started. The controller performs a series of computations with the datapath to generate memory addresses to load the serializer and timing for the print engine, until the page is complete. When the page is complete, it restarts with a new address if the scan pending bit of the command register has been reset.

The mode register, is a read/write. The mode register contents are:

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	Mode Bit:	Reset Level:	Function:
	<0> TMDE	0	Timing mode
	<1> JBMS	0	JBMS mode
5	<2> BBE	0	Band buffer enable
	<3> CINT	0	Print complete interrupt enable
	<4> VINT	0	Vertical interrupt enable
10	<5> PRDY	0	PRDY interrupt enable
	<6> VSREQ	0	VSREQ interrupt enable
	<7> BBI	0	Band buffer interrupt enable
	<8> DIV	0	Video Clock Division select
15	<9> BLV	0	Black Level
	<31:10> reserved	0	Program to 0's

"Timing mode" selects either Primary, JBMS, or Alternate timing mode:

	JBMS	TMDE	MODE
20	0	0	Primary
	0	1	Alternate
	1	0 or 1	JBMS

Band buffer mode is selected if the BBE bit is a 1. If band buffer mode is not enabled, normal full page mode is selected.

CINT, VINT, PRDY and VSREQ are interrupt enables for the respective interrupt sources. If any of these bits are set to 1, the associated interrupt source is enabled to generate the PVP interrupt. More than one interrupt may be active at one time.

When the PRDY bit is set, if the PRDY pin transitions from inactive to active, an interrupt is posted. When the VSREQ bit is set, if the VSREQ pin transitions from inactive to active, an interrupt is posted.

The mode register must be programmed prior to the start of an operation (see the command register, below). The mode register is not double buffered. Changes to the mode register

while the PVP are active is not allowed and may result in unpredictable behavior.

In primary mode, DIV tells the PVP clock generator to divide the video input clock by 8, if 0, or to divide the input clock by 16, if 1. In alternate or JBMS mode, the PVP does not divide the video input clock.

The black level is 1 if the BLV bit is 0 (the default level), otherwise, the black level is 0. When the video is blanked, its level is the same as the white level.

Reserved bits must be programmed to zero's to avoid potential side effects, since some bits may be used for production test features or future enhancements.

The command register, is a write only register. The contents may be read through the status register. The command register contents are:

Command Bit:	Reset Level:	Function:
<0> STF	0	Start frame
<1> RSET	0	Reset PVP (for diagnostics)
<2> BKRL	0	Print Back Page, Right to Left
<3> ERASE	0	Erase After Printing
<31:4> reserved	0	Program to 0's

The command register is used to initiate PVP operations and to support some testing functions. No operations will be initiated other than through the command register. The command register is double buffered. The first loading of the register enables a frame to be generated by the PVP. If the register is loaded a second time, a new frame will commence after the completion of the first, unless the second is a RSET or a testing function. If, subsequently, the command register is loaded a third time, prior to the completion of the first frame, the results may be unpredictable unless the RSET bit is set, in which case all operations will be aborted.

The command which is most important is the "start frame" command, which initiates the timing and DMA operations which the PVP performs during normal system operations. The command register is cleared and all operations halted if the PPC reset pin is asserted. All operations may also be abruptly halted by setting the "RSET" bit of the command register. The RSET bit overrides the other bits.

BKRL indicates that back side printing is to be performed, if it is 1. In this case, the page is scanned in reverse order (from bottom to top and from right to left). ERASE indicates that the frame buffer is to be erased while printing. After a memory word is read for printing, it is erased. Reserved bits must be programmed to zero's to avoid potential side effects, since some bits may be used for production test features or future enhancements.

The status register holds pertinent information for the control and monitoring of the PVP. The bits are unique and have different characteristics. The status register is a read/write register whose contents are:

	Status Bit:	Read/write:	Function:
	<0> FIP	Read	Frame in progress
	<1> FPD	Read	Frame pending
	<2> CINTA	Read/clear	Page completed
	<3> VINTA	Read/clear	Vertical timer completed
25	<4> BBIA	Read/clear	(Band) buffer completed
	<5> PRDYAP	Read/clear	PRDY positive edge transition received
	<6> PRDYAN	Read/clear	PRDY negative edge transition received
30	<7> VSREQA	Read/clear	VSREQ transition received
	<8> LSYNC	Read	LSYNC pin level
	<9> FSYNC	Read	FSYNC pin level
	<10> PRDY	Read	PRDY pin level
	<11> PPRDY	Read	PPRDY pin level
35	<12> Prnt	Read/Write	Print pin level
	<13> CPrdy	Read/Write	CPrdy pin level
	<14> MREQ	Read	Memory request pending
	<15> BLANK	Read	Horizontal blanking status
	<16> FIFOE	Read	FIFO Empty

<17>	FIFOEB	Read	FIFO Block Empty
<18>	FIFO F	Read	FIFO Full
<19>	FIFO RE	Read/clear	FIFO Refill Error
<20>	ERASE	Read	Erase Frame Buffer After
5			Print Enabled
<21>	BKRL	Read	Print Back Page, Right to
			Left
<22>	VSynC	Read/Write	VSynC pin level

10 The FIP bit indicates that a frame has been started, but has not yet completed. The FPD bit indicates that a second operation has been loaded, but has not yet started, except in the case of a reset command, which will halt all operations and clear these bits.

15 The state transition bits (CINTA, VINTA, BBIA, PRDYA, VSREQA) indicate that their associated event has occurred, and that they have not yet been cleared. If their associated interrupt enable has been set, they generate interrupts when active. The interrupts are cleared when the status register state bits are cleared by writing 1's to the bits to be
20 cleared. Writing 0's to the associated bits leaves them unchanged. Once set by the hardware, the associated state bit is not reset except as described above, when the reset pin of the PPC is active, or the PVP reset command is executed.

25 The pin level bits hold the levels of the I/O pins. The LSYNC, FSYNC, PRDY and PPRDY pins are all inputs (therefore read only), while the VSYNC, Prnt and CPRDY pin are outputs. The VSYNC, Prnt and CPrdy pins are cleared to 0 by reset operations, and are programmable to any level by processor writes to the status register.

30 The TMDE, BBE and interrupt enable bits hold the values last programmed into the mode register.

The MREQ status bit indicates that the PVP has one or more outstanding memory requests which have not yet been completed. The BLANK bit indicates that the PVP timing indicates that
35 horizontal blanking is currently occurring. The FIFO status

bits hold various FIFO control states. All of these bits are intended for use in testing.

All interrupts are reset when the PPC is initialized by the reset pin, a reset instruction is set to the PVP command register or the individual bits are reset through the status register.

When the V_int register is loaded, it is decremented on each line reset (when X_current is loaded from X_max). When V_int becomes less than 0, the VINTA status bit is set. An interrupt is sent to the IPU if the VINT interrupt enable bit is set. V_int is generated after the data in the scan line pointed to has been sent to the printer.

The print complete state (CINT) is set after the last scan of data has been sent to the printer. The print complete interrupt is posted if the status bit is set and its interrupt enable is set.

The PRDY interrupt is set when the PRDY pin transitions from inactive to active (PRDYAP transition bit is set in status register) if the PRDY interrupt is enabled. The PRDY interrupt is also set when the PRDY pin transitions from active to inactive (PRDYAN transition bit is set in status register) if the PRDY interrupt is enabled.

The BBI interrupt is set when the PVP is in band buffer mode, the BBI interrupt is enabled, the current buffer has been printed (the BBI status bit is set) and the Base registers have been swapped. When band buffer mode is not selected, the BBI status bit is always cleared.

Timing-

FIGS. 11-14 illustrate the timing of the PPC video interface as used with a conventional laser printer. Synchronization is performed on all interfaces between the video data clocked logic (clocks derived from the Wclk pin) and system-clocked logic. The timing controller contains circuits which are synchronous to both the Wclk and the internal PPC

clocks. It generates timing and control for the serialization hardware, the sequencer and the timing pins. Although the print engine allows the controller to generate the video data clock (Wclk), the PPC does not generate it, but requires the
5 board or print engine to generate the clock as in the alternate mode. This clock may be asynchronous to the PPC clocks.

The sequence of events which occurs in the video processor is:

1) The video register file parameters are pre-loaded, including the base register. This is normally done during the
10 initialization of the printer, unless the page size, frame buffer address or margins are changed.

2) The PRdy pin is asserted by the printer.

3) The Prnt pin is asserted by the controlling software.

15 4) A print command is loaded by the IPU into the command register.

5) When the FSync input is asserted, the sequencer begins to time the top margin (programmed by register T_ margin).

20 6) When the top margin timer completes, the sequencer waits for the next LSync input.

7) When the next LSync input is received, the sequencer begins the left margin timer (programmed by register L_ margin) and fetches the first frame buffer data to be sent to the
25 printer, when the printer is ready.

8) When the left margin timer completes, video data begins to pass out of the PPC chip on the Wdata pin synchronous to the Wclk input clock. The end of the scan is determined, again by the L_ margin counter but using the X_ MAX register data.
30

9) At the end of the scan, the sequencer again waits for the LSync and L_ margin timer until the last scan is completed.

10) When the last scan is complete, the sequencer awaits the next print command. If a pending print command is present,
35 a new frame is immediately started.

The T_margin (FIG. 12) and L margin (FIG. 11) timers are used to synchronize the video data to the paper and the printer scanner, respectively. The timers are synchronized to the FSync and the LSync input pins, respectively.

5 The T_margin counter holds the number of LSync pulses to be counted after FSync but before the start of printing. T_margin timing is performed by the sequencer on the register file. The LSync pulses are synchronized to the PPC clocks and counted. Only the leading edge of Lsync is significant.

10 The L_margin counter initially holds the number of printer video clocks in each line prior to the start of printer data on each scan line. This number is loaded into a hardware counter in the Timing Controller at the start of each L_sync, which is then decremented by Wclk. When the count is
15 exhausted, the video data begins to serialize out of the PPC. The L_margin counter is then re-initialized to the number of pixels within the scan line to be scan converted and used to blank the video at the end of the scan line and to control the reading of the FIFO. When the L_margin count expires, the
20 video data serialization is complete. The VDO is blanked after the expiration of the scan line pixel count.

 Primary mode operation uses UART 0 for interfacing with the printer engine. UART 0 is used to generate serial control signals and receive the serial status information.

25 The operation of UART 0 changes it into a synchronous type of receiver/transmitter. The clock is generated by the PPC internally and passes out of the chip on the U0clk pin. The timing meets the requirements of the Canon print engine. The receive and transmit data circuitry uses a special clock which
30 only transitions when data is to be transmitted to the engine or status returned by the engine. The U0clk frequency is programmed into the UART 0 baud rate generator. FIGS. 13 and 14 illustrate the Canon mode UART 0 output timing and input timing, respectively.

PPC Memory System-

The PPC memory system is illustrated in FIG. 8 and consists of arrays of memory (e.g. 32 or 39 bits, depending on the data bus width and whether ECC is implemented or not) of memory treated as a group of memory words. The arrays are
5 ordered into 2 banks (i.e. 2 way interleaved). The banks consist of one or more RAM arrays and optional ROM arrays. The ROM and RAM interleave, timing and memory depths are independent. Up to 8 arrays per bank are directly supported
10 by the initial PPC, providing up to 16MB of memory with 256K x 4RAM chips or 64MB of memory with 1M x 4RAM chips.

The banks of memory are interleaved through the use of data multiplexing and buffering circuits. The RAM arrays require timing and control signals (RAS, CAS, WE) to be
15 generated and buffered. The PPC generates the required signals, but external buffers are used to provide sufficient drive for the memories and meet their timing requirements. The row and column addresses are multiplexed out of the PPC on one set of address pins. The addresses require some pipelining
20 outside the PPC to make use of the memory bandwidth available from the RAM chips.

FIG. 9 is a diagram of the data chip which interleaves read and write data from the various banks into the PPC during read cycles and holds write data for the individual banks
25 during write cycles. The PPC connects to a bidirectional data port, while the banks connect to individual bidirectional ports.

FIG. 10 is a diagram of the address chip which is used to pipeline row and column addresses to the memories and buffer
30 the memory strobes (RAS, CAS, etc.) for driving the memories. The memory loads may be high, and require buffering between the PPC pins and the high capacitance loads. In addition, the addresses which pass to the memory banks are pipelined differently, depending on which bank is accessed first.

The configuration allows each chip to be used to buffer multiple (e.g. 8) arrays of 32 bits (EEC may be included). One chip is required for small systems which require only a few arrays, but 2 chips may be required for a system which requires 16 arrays. Multiple DRAM organizations such as 256K x 4 or 1M x 4 DRAM's may be used. Using 1M x 1 or 4M x 1 DRAMS will increase the address and control pin loading by approximately a factor of 4. Wider systems, such as 64 bits and other interleaves may be implemented such as 4 way interleaving.

In 64 bit systems, two chips may be used together for each pair of banks resulting in the need for 2 chips for a 2-way interleaved 64 bit system with up to 4 arrays/bank and 4 chips required for a 4-way interleaved 64 bit system. Buffering all memory address and timing strobes with the same chip allows the skew of these signals to be reduced, allowing tighter memory timing to be reliably implemented. Using separate buffer chips increases the skew of the memory timing signals. The configuration of FIG. 9 allows one 128 pin chip to drive 2 banks of 32 bits each containing up to 8 arrays. Other configurations are possible.

I/O Interface-

FIG. 15 is functional block diagram of the I/O controller datapath. The I/O interface controller is designed to provide an interface between the internal IPU, the memory subsystem and external slave and DMA devices. The I/O controller provides and responds to handshaking signals which control the timing and direction of data flow. See copending application SMOS-7992MCF/GBR/WMB. Briefly, the Printer Video Processor (PVP) and the memory Interface are integrated with the I/O controller for sending and receiving addresses and data. It will be noted that the PVP addresses memory through the I/O controller with requests sent through a PVP Request Buffer, but video data is

available to the PVP by direct memory access (DMA) through the PVP Video Data Queue.

Operand alignment logic is used to assemble and align incoming data from external I/O and to re-align outgoing
5 aligned data to external I/O by decoding the lower bits of the address and the byte enables associated with the data. Incoming data from I/O is assembled into aligned blocks, as decoded from the lower bits of the address and placed into an internal data queue. Aligned incoming data from the assembly
10 logic is accumulated in the queue and is sent to memory after a large block of data has been assembled in the queue. The Data is sent to the switch by the queue in a burst, which allows data to be transferred with high bandwidth and minimizes memory subsystem usage by the I/O.

15 Outgoing data from the memory or the IPU, which comes in aligned blocks is re-aligned according to the associated address and data size, and may be passed out of the interface in multiple bus cycles or in burst mode cycles. The number of cycles depends on the data size and the external device size
20 (8 or 16 bits).

Unaligned DMA data transfers can be supported on the I/O bus as long as addresses are contiguous. However, DMA write data and data being read by the IPU must be aligned to the appropriate byte(s) on the bus.

25 IPU reads or writes, which are directed to the PVP, UART's or external I/O devices require that the I/O interface arbitrate for the external I/O bus and perform address decoding before the transactions can be completed. Even though the PVP and UART's are resident on the chip, they are treated almost
30 the same as external I/O devices and access requires the I/O interface to acquire the I/O bus.

The I/O bus is a synchronous 16 bit bus interface. The bus employs separate parallel busses for address and data. The data bus is 16 bits wide, bidirectional. The address bus
35 consists of 31 upper order bit address lines pointing to the

2-byte word location and 2 individual byte enable signals to select active byte within the 2-byte word. The bus is synchronous to the bus clock, which is an output clock generated from the system clock. This clock is intended to be the master clock for all bus decoding and control logic. Dynamic bus sizing is supported by the bus to allow connection of 8 and 16 bit peripherals.

Bus requests by the processor which are more than the bus width (16 bits) or which access devices which are narrower than the data size of the operation will generate multiple bus cycles. An example of the first case is when the processor executes a quad word store to the I/O space, which translates to either eight 16-bit cycles or sixteen 8-bit cycles, depending on the maximum bus size and the *size 8 pins. The I/O bus has a burst mode mechanism to enable high speed data transfers. During burst cycles, a new data items can be clocked out during every clock.

The bus allows external bus masters (DMA devices or other processors), but performs arbitration between the internal I/O controller and external master internally.

External interrupts are also handled by the bus interface dedicated level sensitive interrupt inputs and an external controller are both supported. During the second locked interrupt acknowledge cycle of an external controller interrupt acknowledge sequence, the interrupt vector is read by the interface from the bus and passed to the internal interrupt controller and then to the IPU.

The I/O bus supports non-cacheable bus transfers such as single cycle, multiple cycles, or burst cycles. 8-bit or 16-bit devices are supported. The minimum bus cycle requires two clock cycles. The start of a cycle is signaled by the assertion of the *ADS pin. The end of a bus cycle is indicated by the assertion of the *RDY or *BRDY after *BLAST has become active. A single bus transaction can require multiple data transfers.

The zero wait-state non-burst bus cycle required two clocks. It can support read or write cycles in two clocks (see FIGS. 16 and 17) and is called a 2-2 bus cycle. The assertion of the *ADS at the rising edge of the clock indicates the beginning of a valid bus cycle and that a valid address and bus definition pins are on the address bus and cycle definition lines. The bus cycle is terminated by the combination of *BLAST and *BRDY, or just the assertion of *RDY.

Wait states can be inserted in a bus cycle by driving the *RDY inactive as shown in FIG. 17. By driving *RDY inactive, any number of wait states can be inserted.

Referring to FIG. 18, non-burst mode multiple cycle transfers are performed when the slave asserts RDY (which overrides BRDY). If the master is actually transferring a clock of data, it will hold *BLAST deasserted at the end of every second clock to strobe data into the CPU. If RDY and BRDY are both deasserted, the transfer is extended (wait states are added). Individual transactions may be executed in the same way, except the BLAST is asserted during each individual cycle.

Burst cycles are acknowledged and allowed by the slave if *BRDY is asserted instead of *RDY. *RDY must be driven inactive, as illustrated in FIG. 19. Wait states are inserted into the burst cycle when the slave drives *BRDY and *RDY inactive at the end of second and subsequent clocks. The individual transfers are completed when *BRDY is asserted. Burst mode may be terminated by the slave if it asserts *RDY. The order of burst transfer for data is determined by the first address presented by the IPU. The burst order is one of the cases in the following table:

	First Addr.	Second Addr.	Third Addr.	Fourth Addr.	Fifth Addr.	Sixth Addr.	Seventh Addr.	Eighth Addr.
	0	2	4	6	8	A	C	E
5	8	A	C	E	0	2	4	6

10 The bus allows connection to 8 bit or 16 bit devices using
 the *SIZE 8 pin during bus cycles. The bus size pin is sampled
 on every bus cycle. An 8-bit transfer is selected if *SIZE 8
 is asserted. *SIZE 8 must be asserted before the first
 asserted *RDY or *BRDY is sampled. Multiple bus cycles are
 generated if the device being accessed has smaller bus width
 than the data size of the transfer operation or alignment
 requires multiple cycles. An example is illustrated in FIG.
 15 20.

20 The assertion of the HOLD indicates that another bus
 master is requesting control of the bus. After the current
 bus cycle is completed, the bus is floated (D<31L0>, *BE<1:0>,
 W/R*, D/C*, M/I*, and *ADS are placed in a high impedance
 state). HLDA is asserted to indicate that the processor has
 released the bus to another master. Refer to FIG. 22. The I/O
 bus is relinquished until HOLD is deasserted.

DMA-

25 FIGS. 22 and 23 illustrate the DMA Read cycles and DMA
 Write cycles, respectively.

30 A DMA read cycle starts when ADS is asserted. The ICS pin
 and address are sampled during the rising edge of BClk. The
 data is requested from the MCU. When the data returns, the
 requested data is driven onto the data pins and the IRDY pin
 is asserted. The IRDY pin is asserted to indicate that the
 cycle can be completed, but the cycle does not complete until
 the RDY pin is asserted. BLAST is used to indicate to the I/O

interface that the cycle is the last in a burst. BLAST is always sampled.

A DMA write cycle is similar to a read cycle, except that the I/O is receiving data. It also uses IRDY and BLAST in a similar manner to DMA read cycles. As long as the I/O interface can receive new data, IRDY will remain asserted. If its internal queue fills up, for example, it will deassert IRDY. It also waits for RDY to be asserted before completing the cycle.

For DMA write requests, data and address information is sent to an internal queue. Data is aligned to the location required by the MCU on its way to the queue and once the aligned data is accepted by the queue, the queue will respond by sending a signal back to the DMA device indicating that it is ready to accept more data. If the queue is full, the ready signal is held deasserted. When the external device stops sending data, or when the data can no longer be accumulated by the queue, the MCU is requested to write the data to memory. The MCU will perform the write when its arbiter grants the I/O unit access to the memory port. The data block will then be sent to the switch together with the request information.

For DMA read requests, the address presented by the external DMA device is presented to the MCU. This address is used to address an aligned block of memory (depending on the memory system interleave and bus width). When the data is returned to the I/O queue, the requested data is returned to the requesting DMA device.

PPC Pins-

In the described embodiment of the PPC the following pins are utilized:

Miscellaneous Subtotal: 6

SCLK	<1:0>	in	External Clock Inputs
*Testi		in	Test Input
*Testo		out	Test Output

- 25 -

	*IOts		in	Forces Outputs to Tri-state
	*Sstp		in	IPU Single Step
	Print Engine		9	
5	WData		out	Printer Data Output
	VSynC		out	Vertical Sync Output
	Prnt		out	Printer Print Command Output
	CPrdy		out	Controller Power Ready
	WClk		in	Wdata Serial Clock
10	LSynC		in	Line Sync Strobe
	FSynC		in	Frame Sync Strobe
	PRdy		in	Printer Ready
	PPRdy		in	Printer Power Ready
	Memory Port Subtotal:		72	
15	MData	<31:0>	i/o	Memory Data
	MAdr	<10:0>	out	Row/Column Address
	MRAS*	<1:0>	out	RAS Pins
	MCAS*	<15:0>	out	CAS Pins
	MRRAL*	<1:0>	out	ROM Row Address Latch
	MREN*	<1:0>	out	ROM Enable
20	MWE*	<1:0>	out	Write Enables
	MRD		out	Read
	MODD*		out	Odd Bank Latch Enable
	MOVLD*		out	Odd Valid
	MBK		out	Starting Bank
25	TClk		out	Tracking Clock Output
	I/O Bus Subtotal:		73	
	Pins:	Bits:	I/O:	Functions
	BClk		out	Bus Clock
30	BClkf		in	Bus Clock Frequency/Phase
	Reset		in	Power-up Reset
	ADR	<31:1>	i/o	Address
	*ICS		in	Chip Select
	*BE	<1:0>	i/o	Byte Enables
35	DAT	<15:0>	i/o	Data
	*ADS		i/o	Address Strobe
	*RDY		i	Slave Ready
	*BRDY		i	Burst Ready
	*IRDY		o	Ready Output
	*BLAST		i/o	Burst Last
40	DC		i/o	Bus Cycle Definition
	W/R		i/o	Bus Cycle Definition
	M/IO		i/o	Bus Cycle Definition
	*SIZE8		in	8 Bit Device
45	BREQ		out	Bus Request
	HOLD		in	Bus Acknowledge

	HLDA	out	Bus Grant
	*BOFF	in	Bus Float Pin
	NMI	in	Non-maskable Interrupt
	INTR	in	Maskable interrupt request
5	DINT	in	Direct Maskable Interrupt Inputs
	<3:0>		

Print Engine Control and UART 9

	UOrxd	in	U0 receive data
	*UOcts	in	U0 clear to send
10	UOtxd	out	U0 transmit data
	*UOdtr	out	U0 data terminal ready
	*UOclk	out	U0 Canon CCLK signal
	U1rxd	in	U1 receive data
	*U1cts	in	U1 clear to send
15	U1txd	out	U1 transmit data
	*U1dtr	out	U1 data terminal ready

Summary-

There has been described a page printer controller which utilizes a RISC microprocessor architecture especially designed for control of a page printer. While the invention has been described with reference to a specific embodiment, the description is illustrative of the invention and is not to be construed as limiting the invention. Various modifications and applications may occur to those skilled in the art without departing from the true spirit and scope of the invention as defined by the appended claims.

CLAIMS

1. A page printer controller comprising
an instruction processor,
a printer video processor having a video port,
a data and instruction cache memory,
5 a memory controller for interconnecting said printer video
processor with external memory and with said data and
instruction cache memory, and

an I/O controller for interconnecting said page printer
controller to an I/O bus, said printer video processor being
10 integrated with said I/O controller and said memory interface
for optimization of memory bandwidth and allowing priority
memory access,

said instruction processor interfacing with a host
computer through said I/O bus and including pixel modification
15 circuitry, orthogonal rotator circuitry, and bit/byte mirror
circuitry, and

said printer video processor interfacing with said
instruction processor through said I/O bus and interfacing with
memory through said memory interface and controlling a printer
20 through said video port.

2. The page printer controller as defined by claim 1
wherein all components thereof are formed in a single
semiconductor chip as a very large scale integrated circuit.

25 3. The page printer controller as defined by claim 1
wherein said pixel modification circuitry, said orthogonal
rotator circuitry, and said bit/byte mirror are interconnected
in said instruction processor as internal function units, said
30 instruction processor including a plurality of register files
and multiplex means interconnecting said plurality of register
files with said function units.

4. The page printer controller as defined by claim 1 wherein said printer video processor cooperatively functions with said instruction processor in generating timing signals, performing scan data fetches and serialization, printer contact and printer synchronization.

5. The page printer controller as defined by claim 4 wherein said video processor includes a register file interconnected with said I/O bus and said memory controller, an adder for performing arithmetic operations with said register file and constants for use in generating addresses and control, and a video serializer for receiving parallel data from memory and serializing said data for transfer to a printer through said video port.

6. The page printer controller as defined by claim 5 wherein said printer video processor further includes means for direct memory access whereby video data is received directly from said memory controller.

7. The page printer controller as defined by claim 1 wherein said printer video processor further includes means for direct memory access whereby video data is received directly from said memory controller.

8. A single chip very large scale integrated page printer controller comprising

an instruction processor,

a printer video processor having a video port,

data and instruction cache memory,

a memory controller for interconnecting said printer video processor with external memory and with said data and instruction cache memory, and

an I/O controller for interconnecting said page printer controller to an I/O bus, said printer video processor being

integrated with said I/O controller and with data directly supplied by a memory interface unit for optimization of memory bandwidth and allowing priority memory access,

5 said instruction processor interfacing with a host computer through said I/O bus, said printer video processor interfacing with said instruction processor through said I/O bus and interfacing with memory through said memory interface and controlling a printer through said video port.

10 9. The page printer controller as defined by claim 8 wherein said printer video processor further includes means for direct memory access whereby video data is received directly from said memory controller unit.

15 10. The page printer controller as defined by claim 9 wherein said printer video processor cooperatively functions with said instruction processor in generating timing signals, performing scan data fetches and serialization, and printer synchronization.

20 11. The page printer controller as defined by claim 10 wherein said video processor includes a register file interconnected with said I/O bus and said memory controller, an adder for performing arithmetic operations with said register file and constants for use in generating addresses and control, and a video serializer for receiving parallel data from memory and serializing said data for transfer to a printer
25 though said video port.

30 12. The page printer controller as defined by claim 8 wherein said video processor includes a register file interconnected with said I/O bus and said memory controller, an adder for performing arithmetic operations with said register file and constants for use in generating addresses and control, and a video serializer for receiving parallel data

from memory and serializing said data for transfer to a printer though said video port.

13. The page printer controller as defined by claim 8 wherein said instruction processor is a RISC processor including pixel modification circuitry, orthogonal rotator circuitry, and bit/byte mirror circuitry integrated therein as operating as integral function units within said RISC processor.

14. A page printer controller comprising
an instruction processor,
a printer video processor having a video port,
a data and instruction cache memory,
a memory controller for interconnecting said printer video processor with external memory and with said data and instruction cache memory, and

an I/O controller for interconnecting said page printer controller to an I/O bus, said printer video processor being integrated with said I/O controller and said memory interface for optimization of memory bandwidth and allowing priority memory access,

said instruction processor interfacing with a host computer through said I/O bus, and

said printer video processor interfacing with said instruction processor through said I/O bus and interfacing with memory through said memory interface and controlling a printer through said video port.

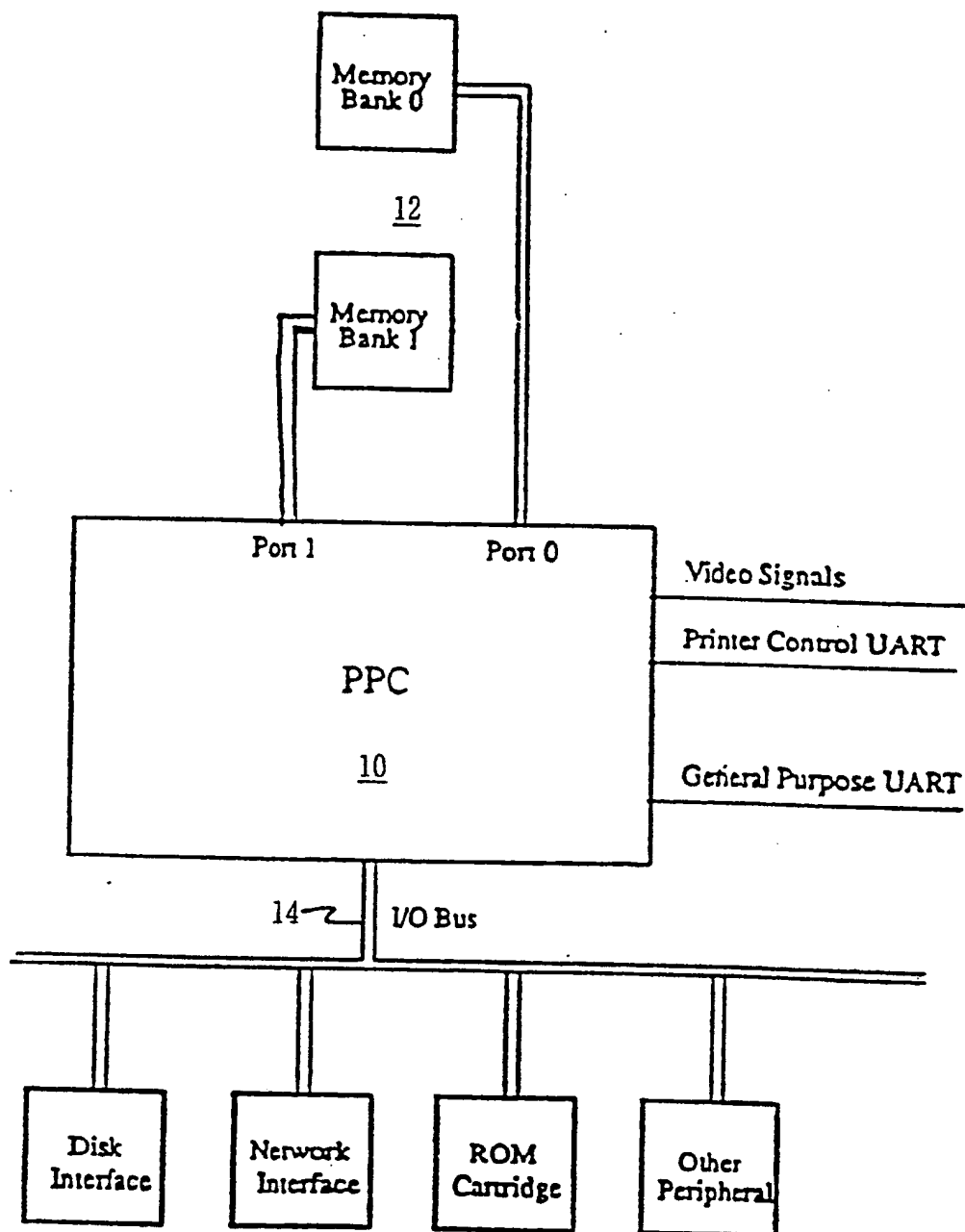


FIG. 1

PPC Block Diagram

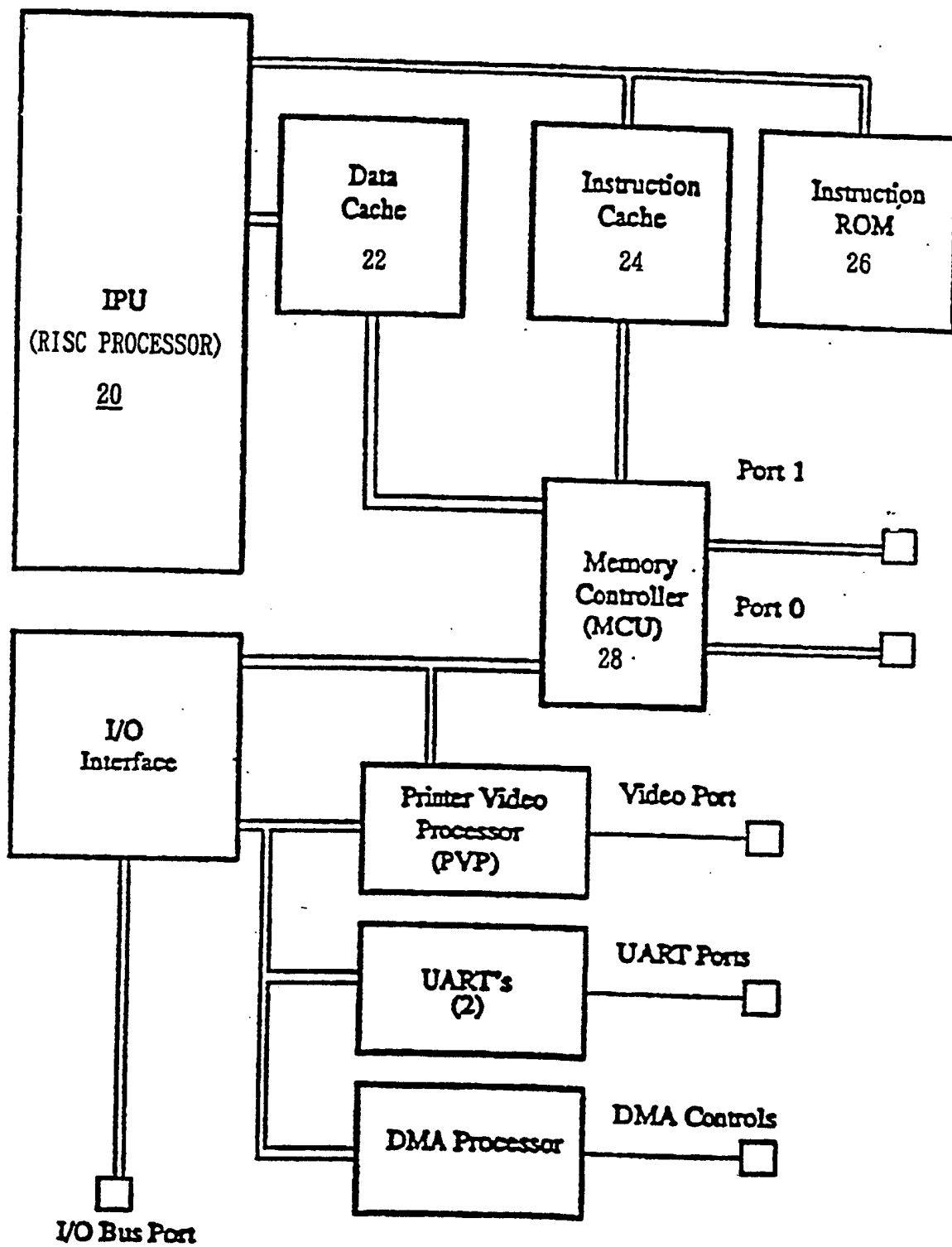


FIG. 2

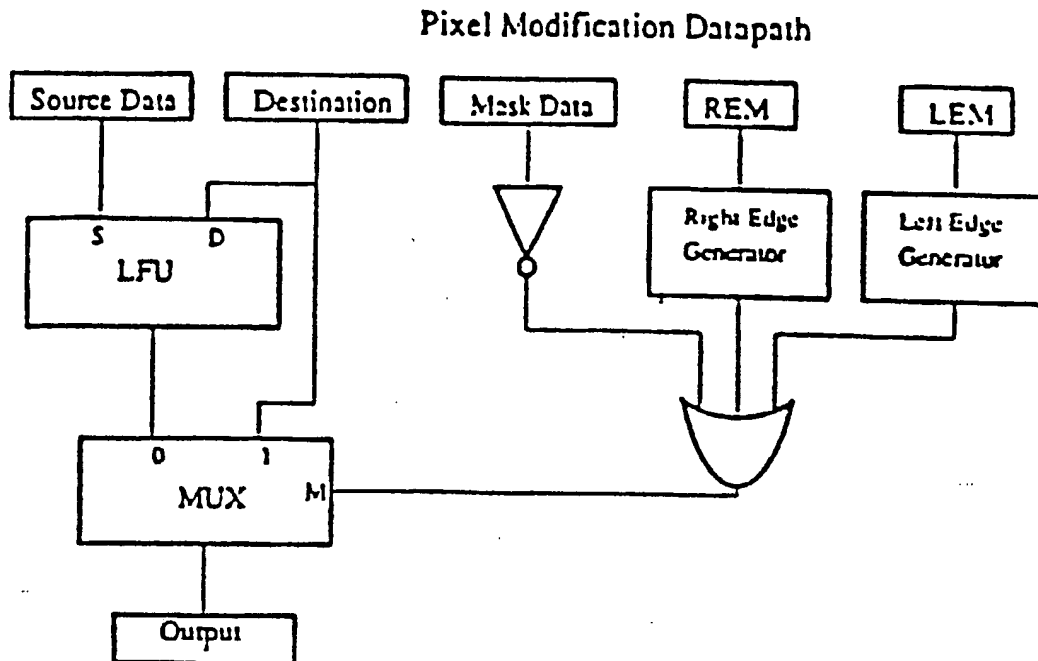


FIG. 3

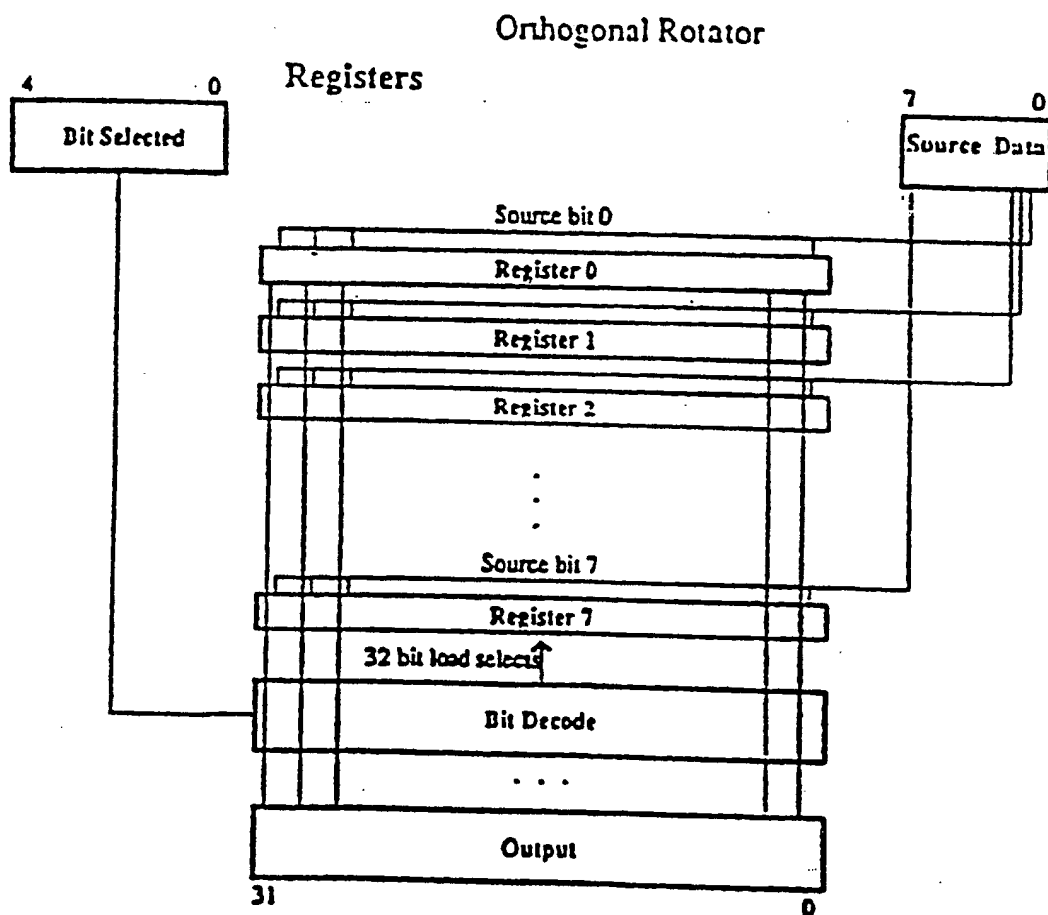


FIG. 4

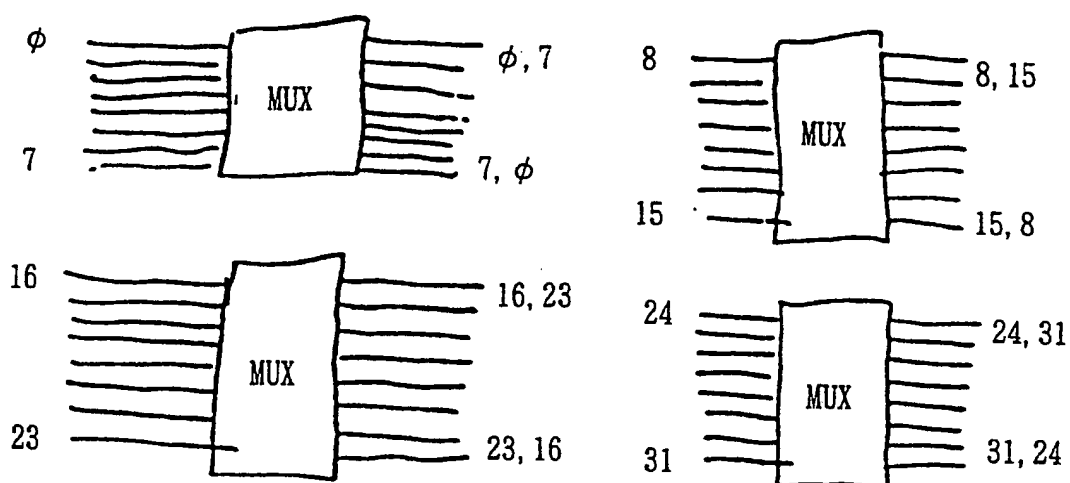


FIG. 5

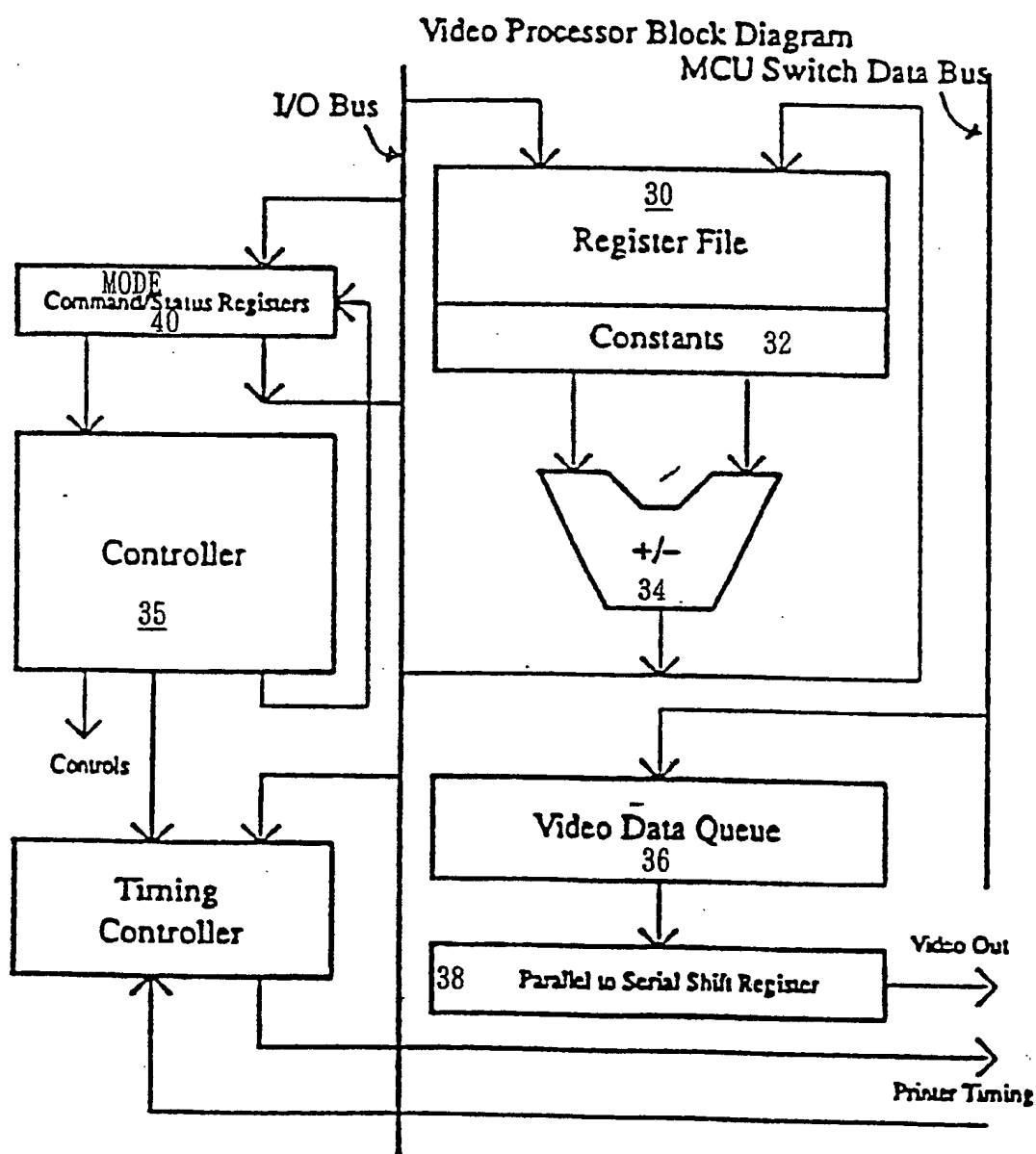


FIG. 7

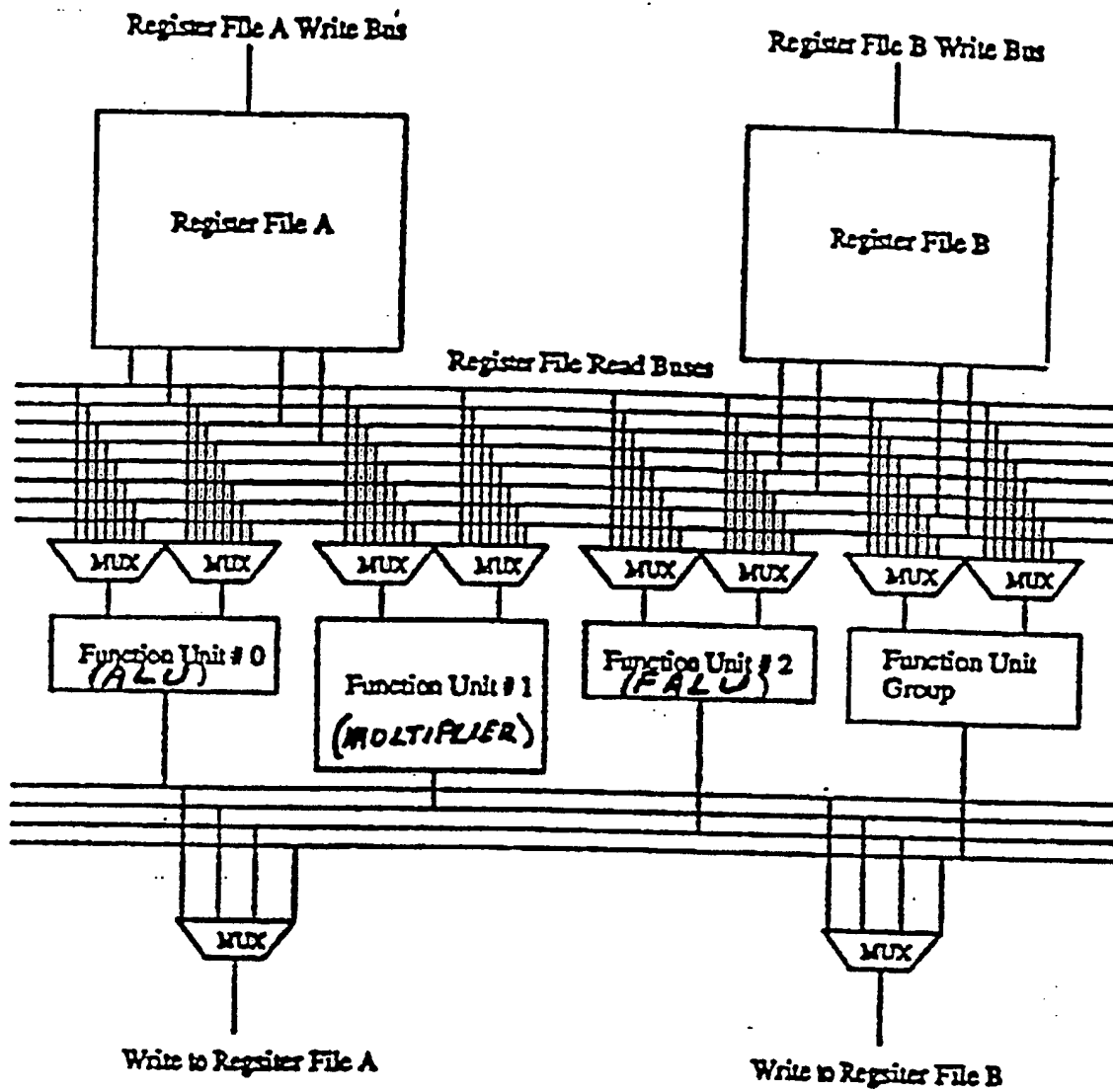


FIG. 6A

6/16

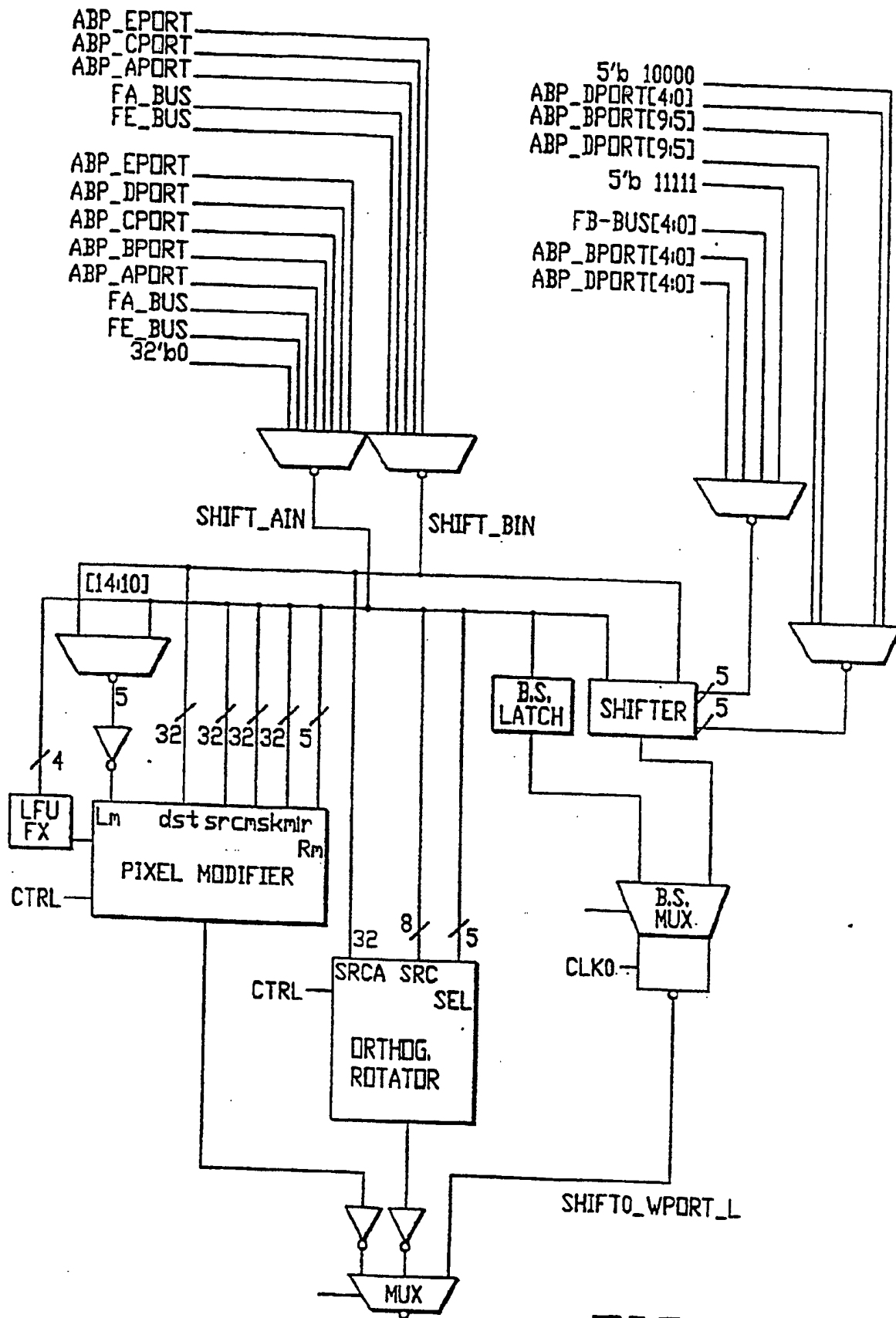


FIG.-6B

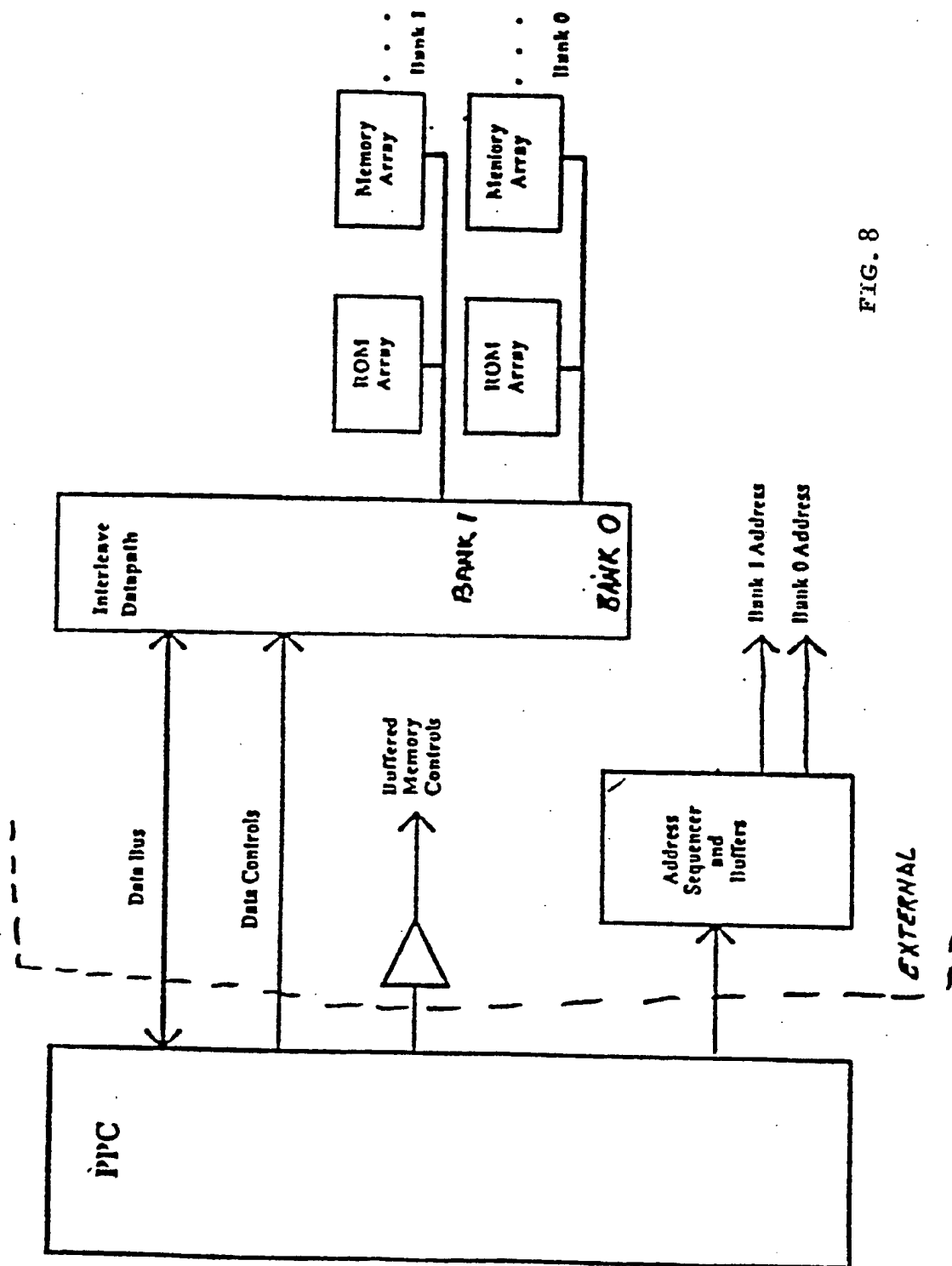


FIG. 8

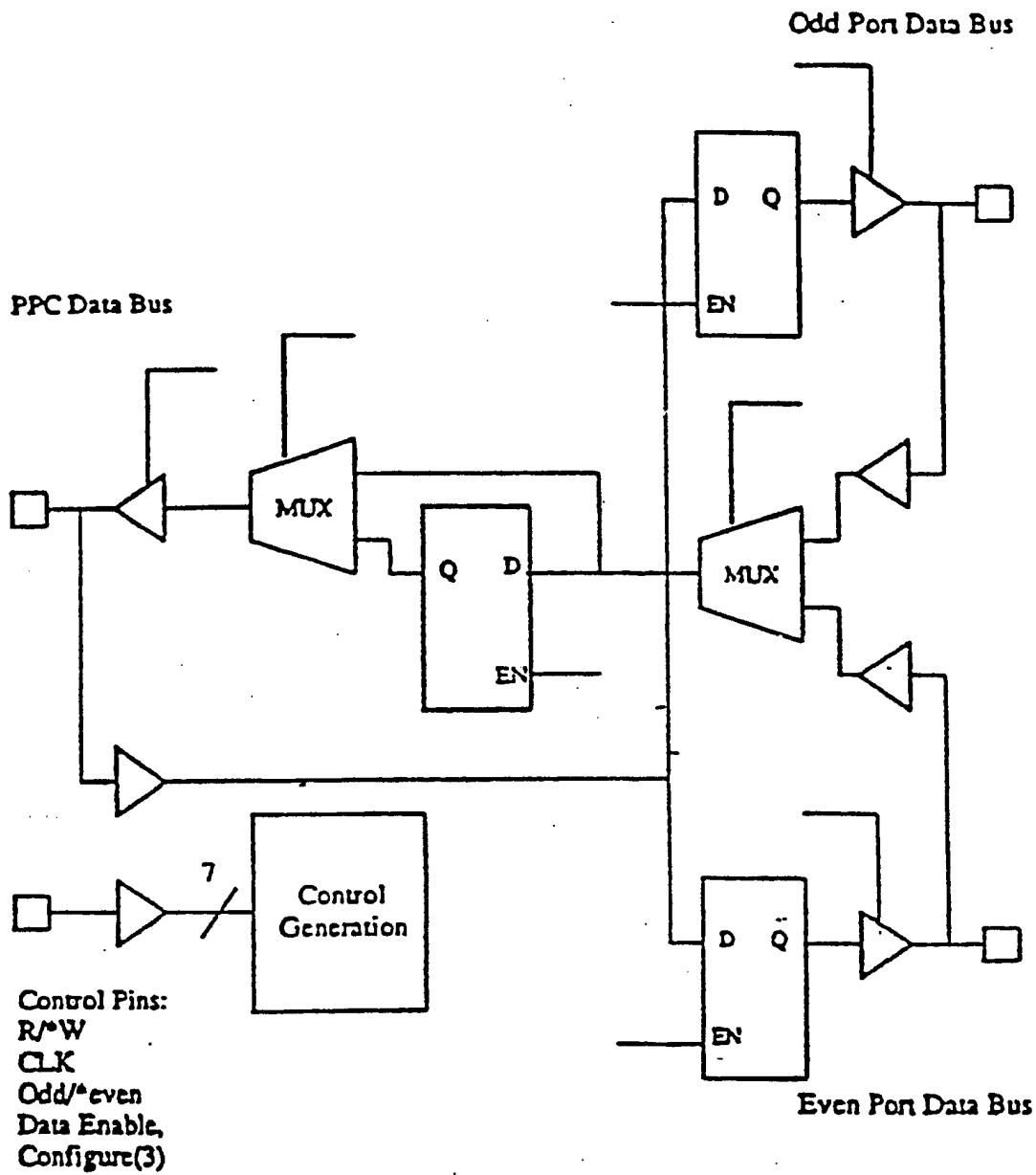
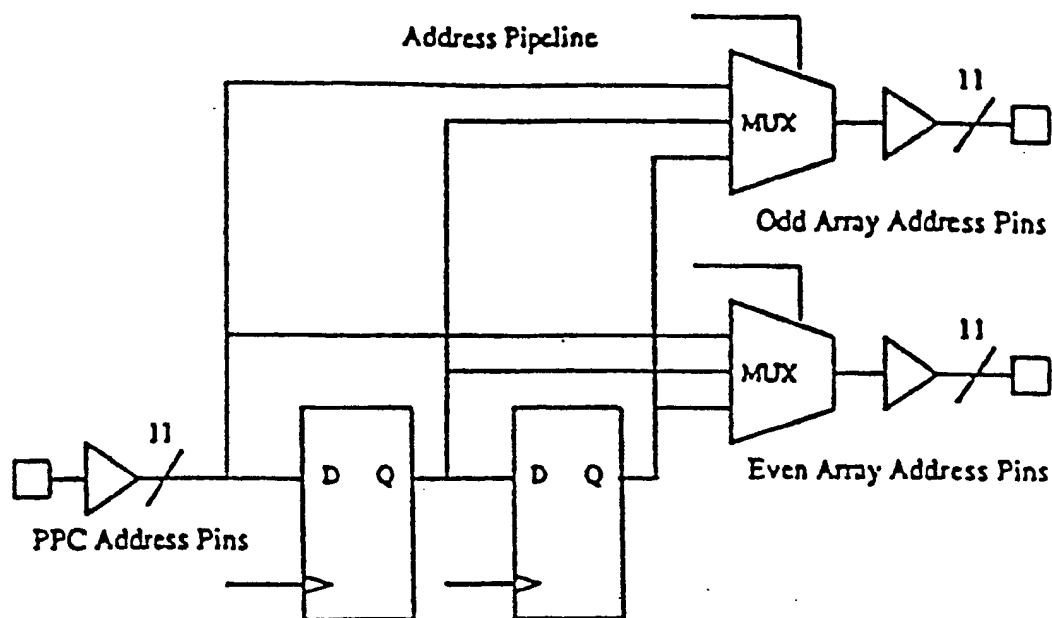
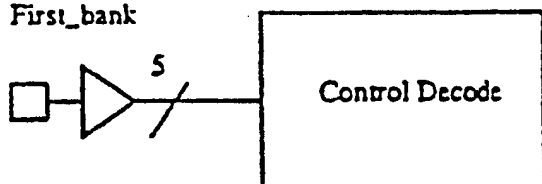


FIG. 9

Address Chip, 2-way Interleave



Control Inputs:
Configuration,
Clock,
Read,
First_bank



Strobe Inputs:
ROM_EN<1:0>,
RAS<1:0>,
CAS<7:0>,
WE<7:0>

Strobe Outputs:
ROM_EN<1:0>,
RAS<1:0>,
CAS<7:0>,
WE<1:0>

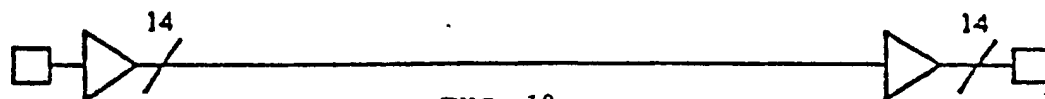


FIG. 10

Horizontal Timing

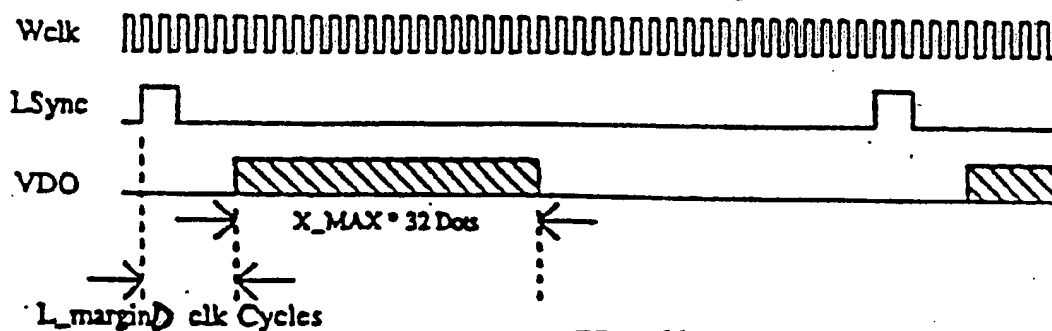


FIG. 11

10/16

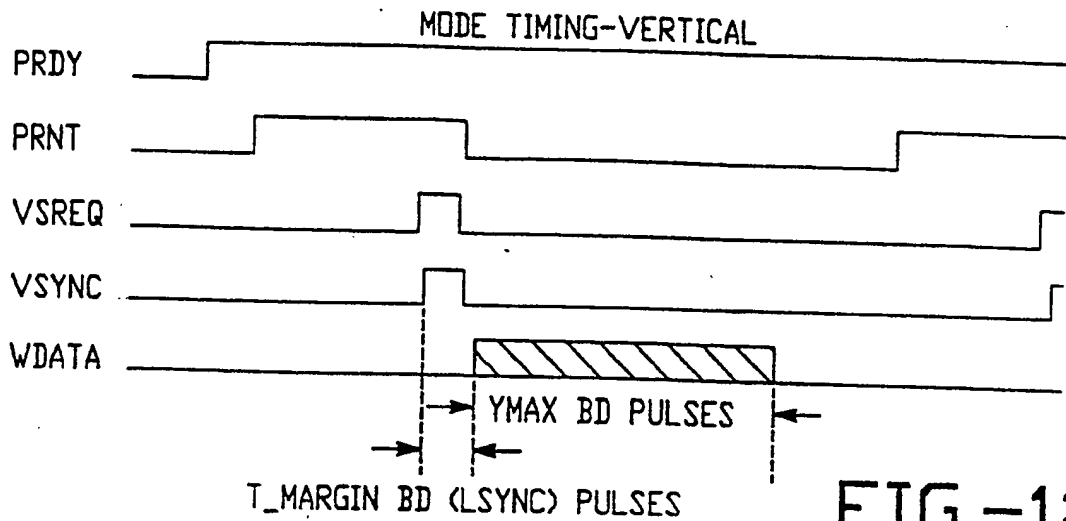


FIG.-12

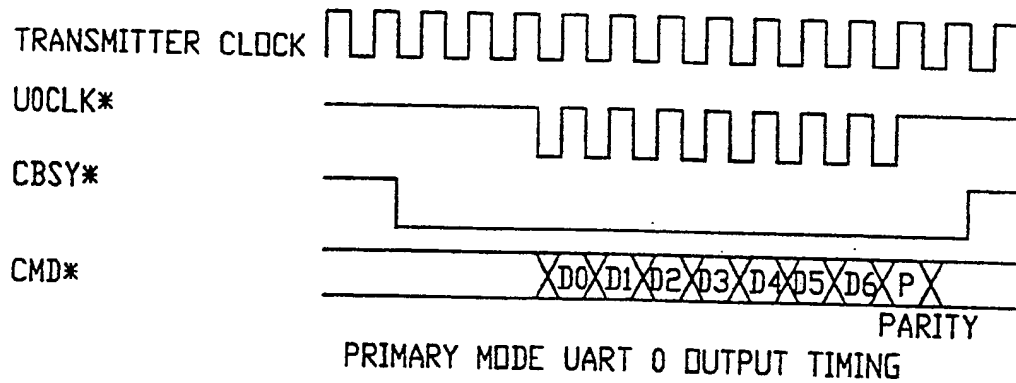


FIG.-13

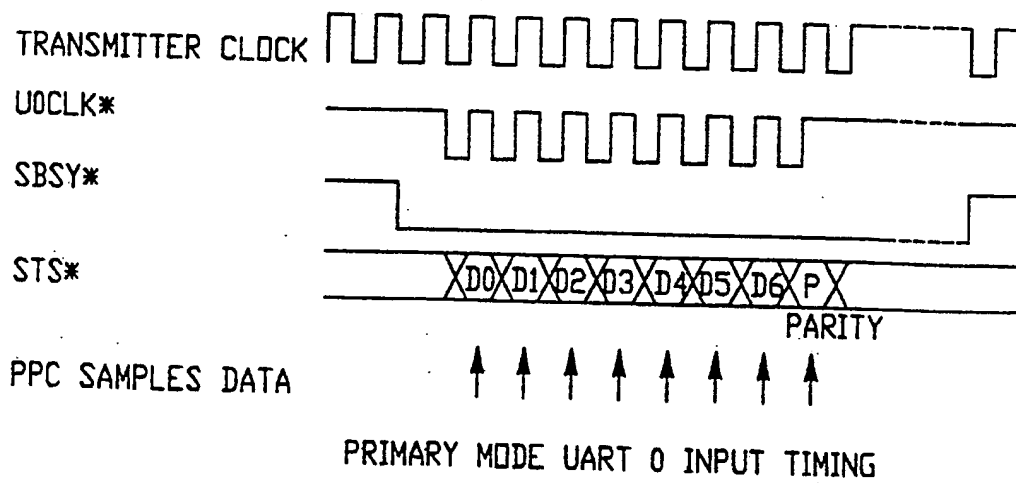
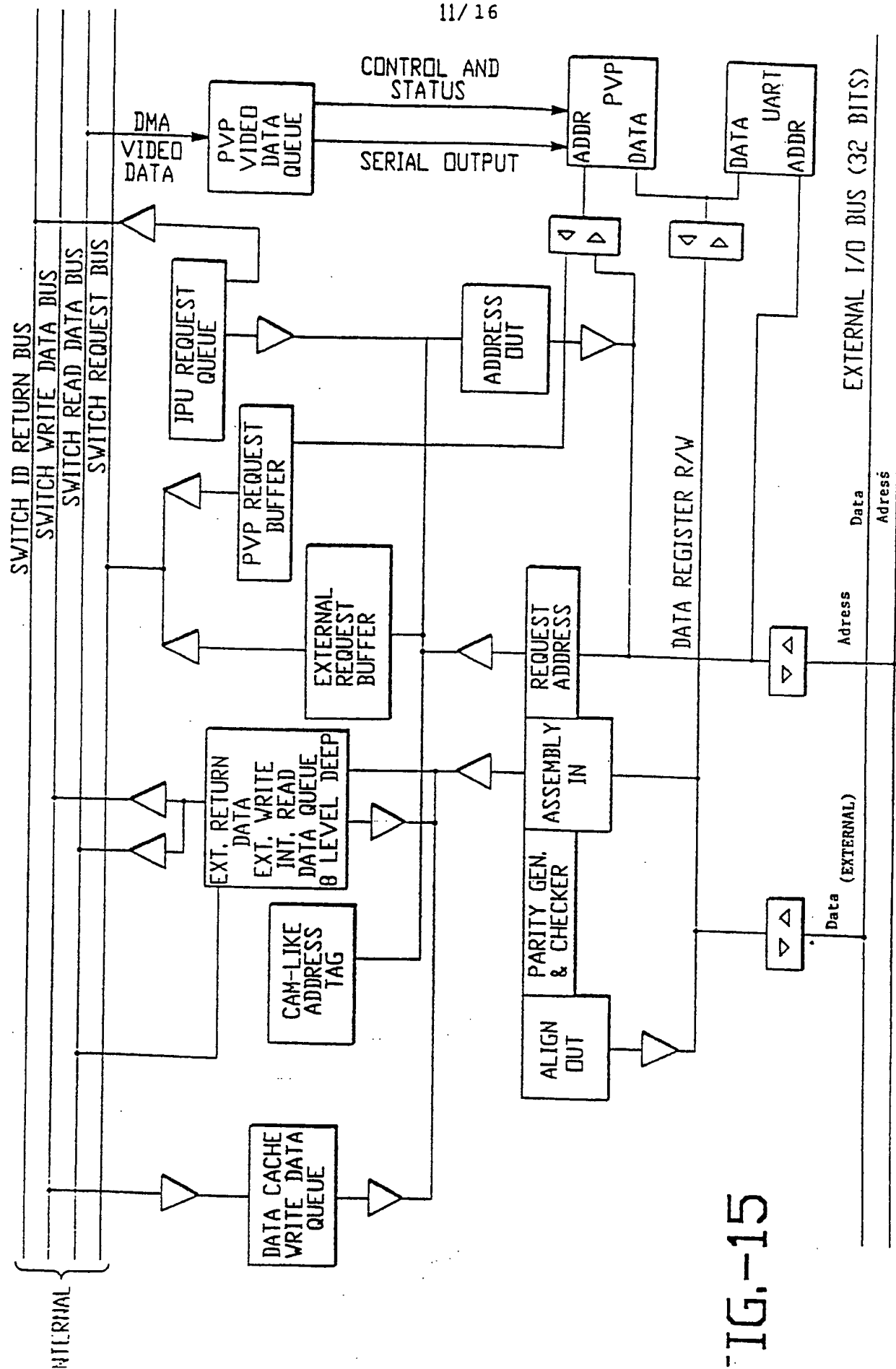


FIG.-14



15-15

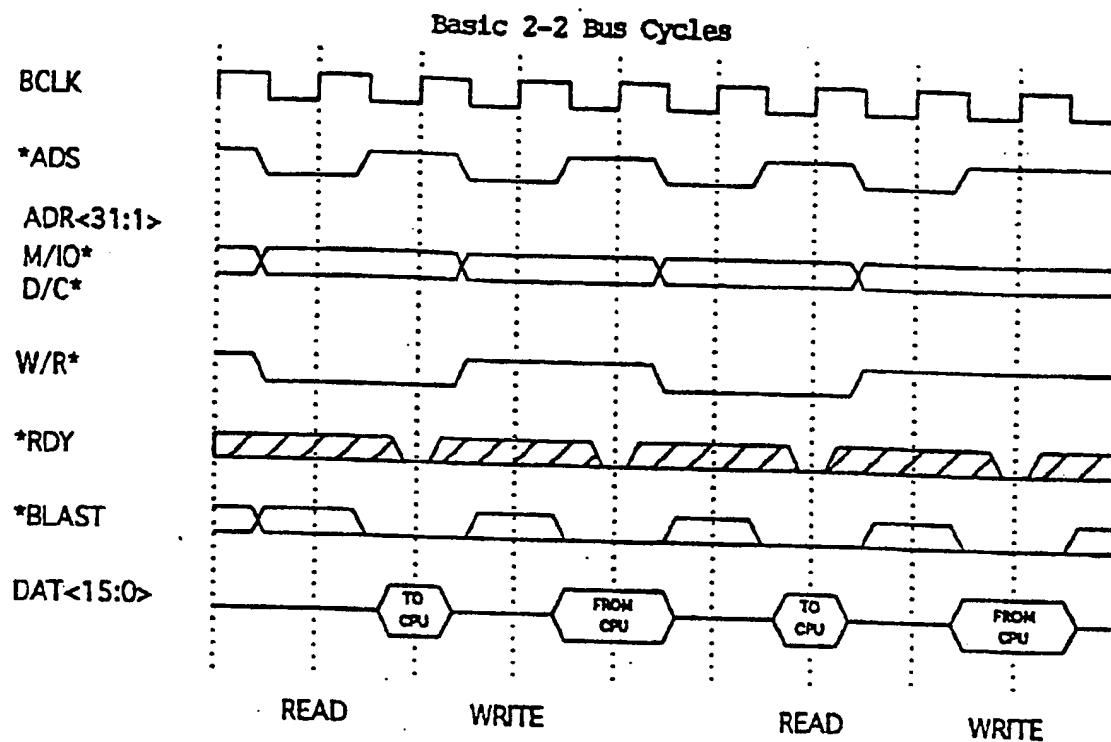
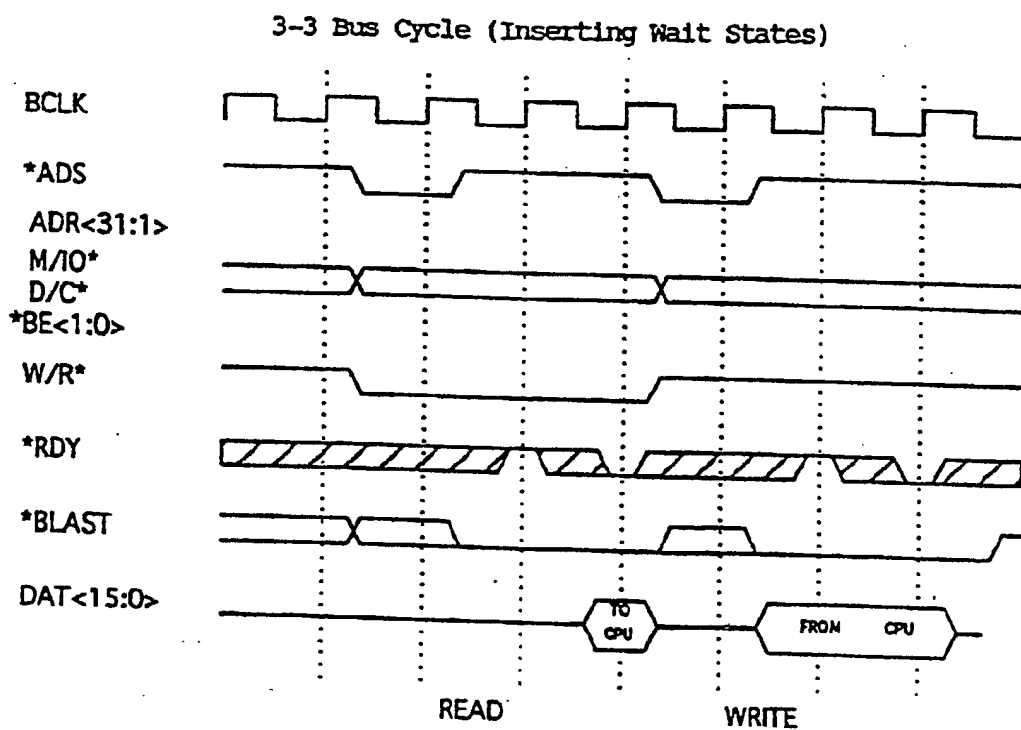


FIG.-16



Non-Burst, Multiple Cycle Transfers

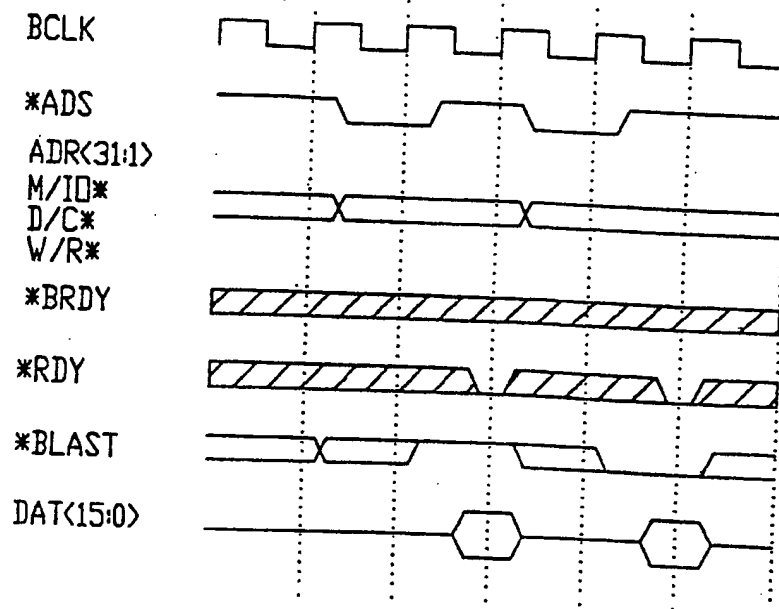


FIG.-18

Burst Mode Cycles

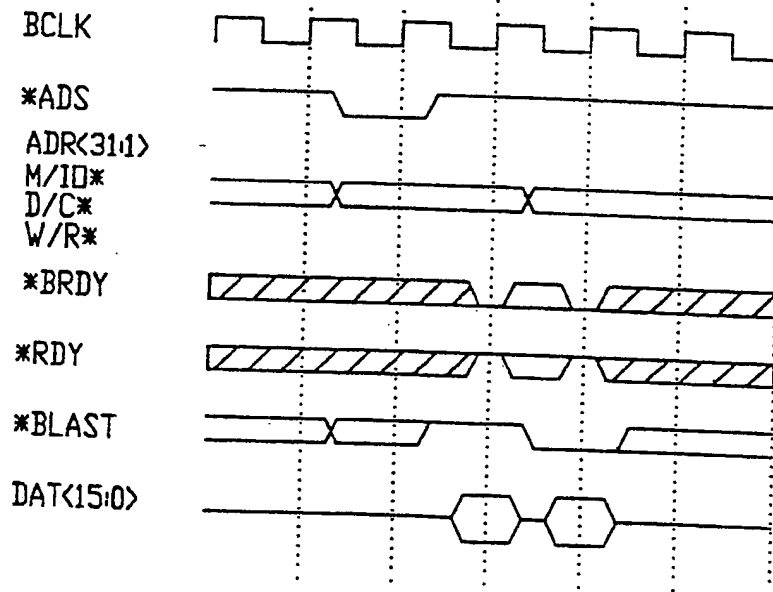


FIG.-19

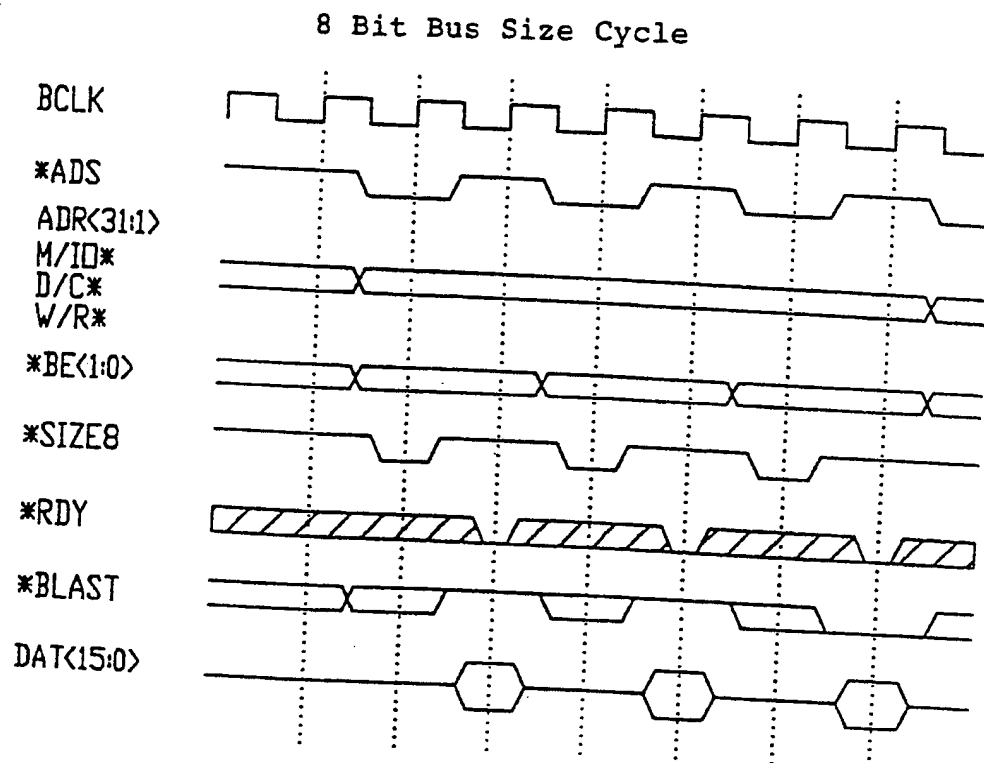


FIG.-20

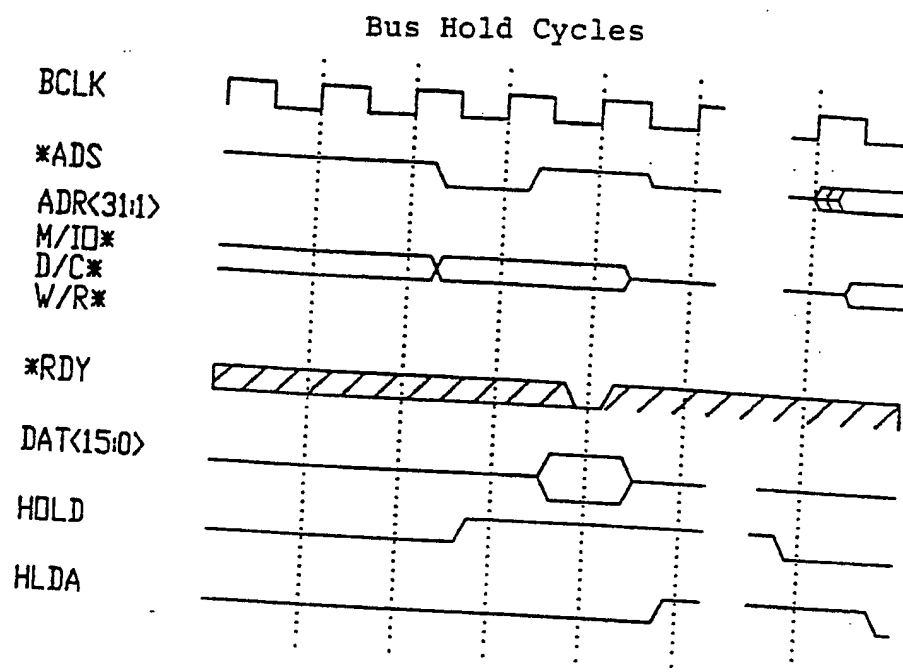


FIG.-21

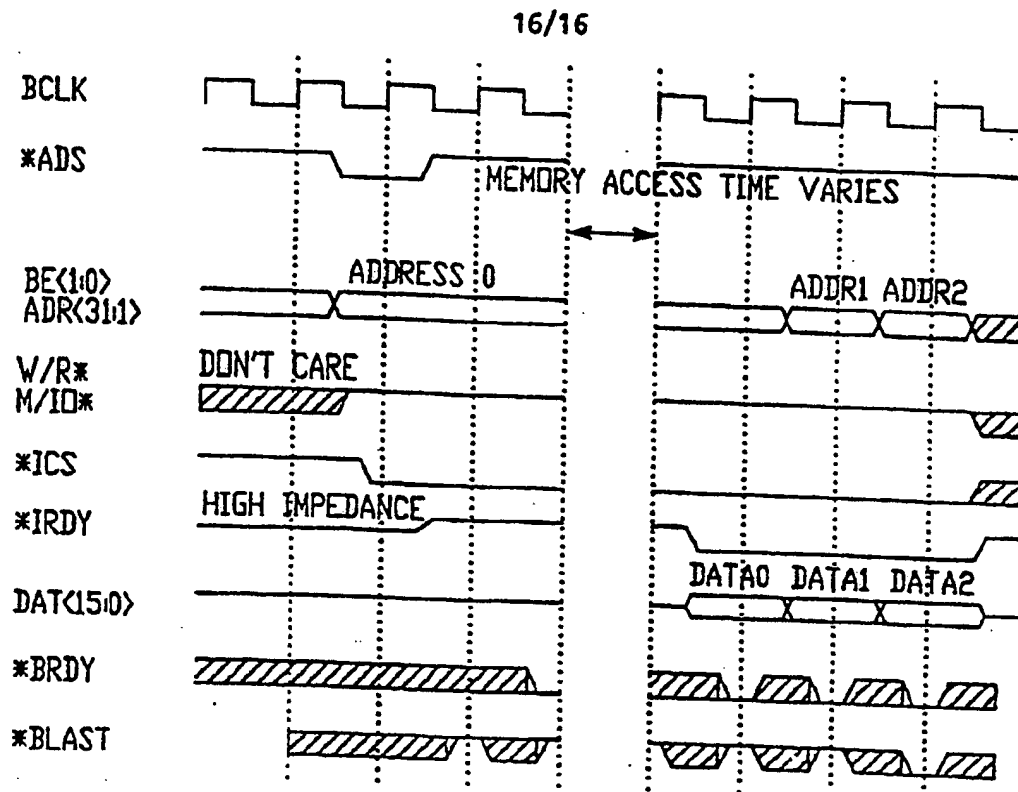


FIG.-22

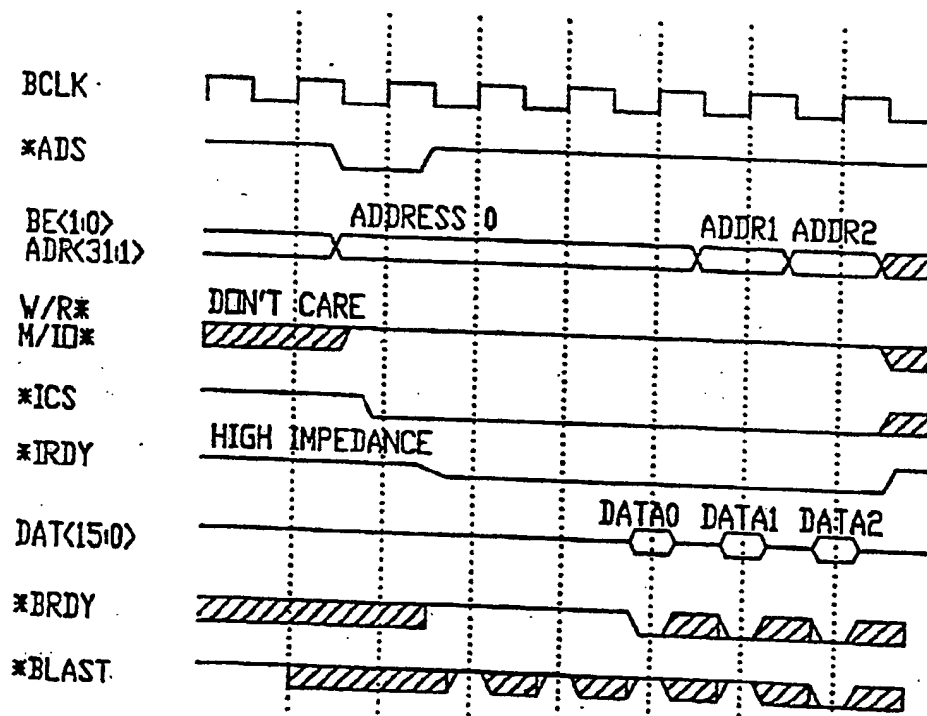


FIG.-23

INTERNATIONAL SEARCH REPORT

International Application No

PCT/JP 92/00867

I. CLASSIFICATION OF SUBJECT MATTER (if several classification symbols apply, indicate all)⁶

According to International Patent Classification (IPC) or to both National Classification and IPC

Int.Cl. 5 G06K15/00; G06K15/02

II. FIELDS SEARCHEDMinimum Documentation Searched⁷

Classification System

Classification Symbols

Int.Cl. 5

G06K

Documentation Searched other than Minimum Documentation
to the Extent that such Documents are Included in the Fields Searched⁸**III. DOCUMENTS CONSIDERED TO BE RELEVANT⁹**

Category ¹⁰	Citation of Document, ¹¹ with indication, where appropriate, of the relevant passages ¹²	Relevant to Claim No. ¹³
A	EP,A,0 354 024 (HITACHI LTD., TOKYO, JP) 7 February 1990 see column 1, line 9 - line 13 see column 1, line 43 see column 4, line 12 - line 17 ---	1,8
A	IBM TECHNICAL DISCLOSURE BULLETIN vol. 33, no. 8, January 1991, ARMONK, NY, USA pages 485 - 486 'Printer using direct memory access to reduce communications bandwidth' ---	6,7,9
A	US,A,4 648 045 (DEMETRESCU) 3 March 1987 -----	2,8

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Date of the Actual Completion of the International Search

30 OCTOBER 1992

Date of Mailing of this International Search Report

06.11.92

International Searching Authority

EUROPEAN PATENT OFFICE

Signature of Authorized Officer

VEEN G.E.

ANNEX TO THE INTERNATIONAL SEARCH REPORT
ON INTERNATIONAL PATENT APPLICATION NO.

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This annex lists the patent family members relating to the patent documents cited in the above-mentioned international search report.
The members are as contained in the European Patent Office EDP file on
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